

**Tulsiramji Gaikwad-Patil College of Engineering & Technology**  
**Department of Master in Computer Application**



Subject Notes  
Academic Session: 2018 – 2019

**Subject:DMGT**

**Semester: II**

Either

1.(A) Let A, B and C be finite sets, then prove that :

$$|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|.$$

**Solution**

We have  $A \cup B \cup C = \{a, b, c, d, e, g, h, k, m, n\}$ ,  $A \cap B = \{a, b, e\}$ ,  $A \cap C = \{b, d, e\}$ ,  $B \cap C = \{b, e, g, h\}$ , and  $A \cap B \cap C = \{b, e\}$ , so  $|A| = 5$ ,  $|B| = 5$ ,  $|C| = 8$ ,  $|A \cup B \cup C| = 10$ ,  $|A \cap B| = 3$ ,  $|A \cap C| = 3$ ,  $|B \cap C| = 4$ , and  $|A \cap B \cap C| = 2$ .  
Thus  $|A| + |B| + |C| - |A \cap B| - |B \cap C| - |A \cap C| + |A \cap B \cap C| = 5 + 5 + 8 - 3 - 3 - 4 + 2$   
or 10, and Theorem 3 is verified. ♦

(B) Prove

$$5 + 10 + 15 + \dots + 5n = 5n(n+1)/2$$

for all  $n \geq 1$ ; using Mathematical Induction

Sol==

using mathematical induction for all  $n \geq 1$ .

Sol<sup>n</sup>:  $\Rightarrow$  for  $n = 5$ .

$$\text{L.H.S.} = P(5) = 5 \times 5 = 25$$
$$\text{R.H.S.} = P(5) = \left( \frac{5 \times 1 (1+1)}{2} \right) = \left( \frac{5 \times 2}{2} \right) = \frac{10 \times 5}{2} = 25$$

$\therefore P(5)$  is true for  $n = 5$

L.H.S. = C.P.K

Induction step: For  $n = k$

$$\text{L.H.S.} = P(k) = 5 \times 1 + 5 \times 2 + 5 \times 3 + 5 \times 4 + \dots + k \cdot n$$
$$= \left[ \frac{k \times n (k+1)}{2} \right]$$
$$P(k) : \text{R.H.S.} = \left[ \frac{5 \cdot n (n+1)}{2} \right]$$

$\therefore P(k)$  is true for  $n = k$  - - - (ii)

$$\begin{aligned}
 P(k+1) &= LHS = 5 + 10 + 15 + \dots + k + (k+1) \\
 &= \left[ \frac{k(k+1)}{2} \right] + (k+1) \\
 &= \frac{k^2(k+1) + 4(k+1)(k+1)}{2} \\
 &= \frac{(k+1) [k^2 + 4(k+1)]}{2} \\
 &= \frac{(k+1) [k^2 + 4k + 4]}{2} \\
 &= \frac{(k+1) (k+2)}{2} \\
 &= \left[ \frac{(k+1)(k+2)}{2} \right] \\
 &= R.H.S.
 \end{aligned}$$

(c) Prove that :

If A, B and C are Boolean matrices of compatible sizes then,

$$(A B) C = A (B C)$$

**Sol== Solution:**

Let us assume that

$$A = [a_{ij}] m \times n$$

$$B = [b_{jk}] m \times n$$

$$C = [c_{kl}] m \times n$$

$$(A \vee B) \vee C = A \vee (B \vee C)$$

$$\text{i.e. } (A+B)+C = A+(B+C)$$

Now,

$$A + B = [a_{ij}] m \times n + [b_{jk}] m \times n$$

$$= [a_{ij} + b_{jk}] m \times n$$

$$(A + B) + C = [a_{ij} + b_{jk}] m \times n + [c_{kl}] m \times n$$

$$= [a_{ij} + b_{jk} + c_{kl}] m \times n \longrightarrow (1)$$

$$(B + C) = [b_{jk}] m \times n + [c_{kl}] m \times n$$

$$=[b_{jk} + c_{kl}] m \times n$$

$$\begin{aligned} (B + C) + A &= [a_{ij}] m \times n + [b_{jk} + c_{kl}] m \times n \\ &= [a_{ij} + b_{jk} + c_{kl}] m \times n \quad \longrightarrow \quad (2) \end{aligned}$$

From equation (1) & (2)

We get  $(A \vee B) \vee C = A \vee (B \vee C)$

i.e.

$$(A \oplus B) \oplus C = A \oplus (B \oplus C)$$

(d) Obtain conjunctive normal form of :

$$\bar{u}(P \vee Q) \quad \text{---} \quad (P \wedge Q).$$



**Proof:-**

$$\begin{aligned} \bar{u}(P \vee Q) &\leftrightarrow (P \wedge Q) \\ \text{by } R \leftrightarrow S &\leftrightarrow (R \rightarrow S) \wedge (S \rightarrow R) \\ &\Leftrightarrow [\bar{u}(P \vee Q) \rightarrow (P \wedge Q)] \wedge [(P \wedge Q) \rightarrow \bar{u}(P \vee Q)] \\ &\Leftrightarrow [\bar{u}\bar{u}(P \vee Q) \vee (P \wedge Q)] \wedge [\bar{u}(P \wedge Q) \vee \bar{u}(P \vee Q)] \quad \{P \rightarrow Q \Rightarrow \bar{u}P \vee Q\} \\ &\Leftrightarrow [(P \vee Q) \vee (P \wedge Q)] \wedge [(\bar{u}P \vee \bar{u}Q) \vee \bar{u}(P \vee Q)] \quad \{\text{By Demorgans property} \& \bar{u}\bar{u}P \Rightarrow P\} \\ &\Leftrightarrow [(P \vee Q \vee P) \wedge (P \vee Q \vee Q)] \wedge [(\bar{u}P \vee \bar{u}Q \vee \bar{u}P) \wedge (\bar{u}P \vee \bar{u}Q \vee \bar{u}Q)] \quad \{\text{By Distributive property}\} \\ &\Leftrightarrow ((P \vee P) \vee Q) \wedge ((Q \vee Q) \vee P) \wedge ((\bar{u}P \vee \bar{u}P) \vee \bar{u}Q) \wedge ((\bar{u}Q \vee \bar{u}Q) \vee \bar{u}P) \quad \{\text{By Associative property}\} \\ &\Leftrightarrow (P \vee Q) \wedge (Q \vee P) \wedge (\bar{u}P \vee \bar{u}Q) \wedge (\bar{u}Q \vee \bar{u}P) \quad \{P \vee P = P\} \\ &\Leftrightarrow (P \vee Q) \wedge (P \vee Q) \wedge (\bar{u}P \vee \bar{u}Q) \wedge (\bar{u}P \vee \bar{u}Q) \quad \{\text{By commulative property}\} \\ &\Leftrightarrow (P \vee Q) \wedge (\bar{u}P \vee \bar{u}Q) \quad \{P \wedge P = P\} \end{aligned}$$

It is the form of product of elementary sum of min terms.

Hence it is form of Principal Conjunction Normal Form.

(c) Obtain Principal Disjunctive Normal Form of

$$P \rightarrow ((P \rightarrow Q), \wedge \bar{u}(\bar{u}Q \vee \bar{u}P)).$$

$$\begin{aligned} \text{Sol} &= \Rightarrow (\bar{u}P \wedge T) \vee (Q \wedge \bar{u}T) \quad \{\text{by } P \wedge T = P\} \\ &\Rightarrow [\bar{u}P \wedge (Q \vee \bar{u}Q)] \vee [Q \wedge (P \vee \bar{u}P)] \quad \{\text{by } P \vee \bar{u}P = T\} \end{aligned}$$

$$\begin{aligned} &\Rightarrow [(\neg P \wedge Q) \vee (\neg P \wedge \neg Q)] \vee [(Q \wedge P) \vee (Q \wedge \neg P)] && \text{\{by distributive property\}} \\ &\Rightarrow (\neg P \wedge Q) \vee (\neg P \wedge \neg Q) \vee (Q \wedge P) \vee (Q \wedge \neg P) && \text{\{by Associative property\}} \\ &\Rightarrow (\neg P \wedge Q) \vee (\neg P \wedge \neg Q) \vee (P \wedge Q) \vee (P \wedge \neg Q) && \text{\{by Commutative property\}} \end{aligned}$$

$\therefore$  It is form of sum of elementary product of min term.

Hence, it is in the form of Principal Disjunction Normal Form.

$$\begin{aligned} &\neg(P \vee Q) \leftrightarrow (P \wedge Q) \\ &\text{by } R \leftrightarrow S \leftrightarrow (R \rightarrow S) \wedge (S \rightarrow R) \\ &\Leftrightarrow [\neg(P \vee Q) \rightarrow (P \wedge Q)] \wedge [(P \wedge Q) \rightarrow \neg(P \vee Q)] \\ &\Leftrightarrow [\neg\neg(P \vee Q) \vee (P \wedge Q)] \wedge [\neg(P \wedge Q) \vee \neg(P \vee Q)] && \{P \rightarrow Q \Rightarrow \neg P \vee Q\} \\ &\Leftrightarrow [(P \vee Q) \vee (P \wedge Q)] \wedge [(\neg P \vee \neg Q) \vee \neg(P \vee Q)] && \{\text{By Demorgans property \& } \neg\neg P \Rightarrow P\} \\ &\Leftrightarrow [(P \vee Q \vee P) \wedge (P \vee Q \vee Q)] \wedge [(\neg P \vee \neg Q \vee \neg P) \wedge (\neg P \vee \neg Q \vee \neg Q)] && \{\text{By Distributive property}\} \\ &\Leftrightarrow ((P \vee P) \vee Q) \wedge ((Q \vee Q) \vee P) \wedge ((\neg P \vee \neg P) \vee \neg Q) \wedge ((\neg Q \vee \neg Q) \vee \neg P) && \{\text{By Associative property}\} \\ &\Leftrightarrow (P \vee Q) \wedge (Q \vee P) \wedge (\neg P \vee \neg P) \wedge (\neg Q \vee \neg P) && \{P \vee P = P\} \\ &\Leftrightarrow (P \vee Q) \wedge (P \vee Q) \wedge (\neg P \vee \neg P) \wedge (\neg P \vee \neg Q) && \{\text{By commulative property}\} \\ &\Leftrightarrow (P \vee Q) \wedge (\neg P \vee \neg Q) && \{P \wedge P = P\} \end{aligned}$$

(b) Find an explicit formula for the sequence defined by  $C_n = 6C_{n-1} + 7C_{n-2}$  with initial conditions

$$C_0 = 2, C_1 = 1.$$

Sol== First find sequence for recurrence relation

$$a_n = 4a_{n-1} + 5a_{n-2}$$

$$\text{For } n = 3 \quad a_3 = 4a_{3-1} + 5a_{3-2}$$

$$= 4a_2 + 5a_1$$

$$= 4(6) + 5(2)$$

$$= 24 + 10$$

$$= 34$$

$$\text{For } n = 4 \quad a_4 = 4a_{4-1} + 5a_{4-2}$$

$$= 4a_3 + 5a_2$$

$$= 4(34) + 5(6)$$

$$= 166$$

$$\text{For } n=5 \quad a_5 = 4a_{5-1} + 5a_{5-2}$$

$$\begin{aligned}
&= 4a_4 + 5a_3 \\
&= 4(166) + 5(34) \\
&= 834
\end{aligned}$$

∴ Sequence is 2,6,34,166,834----

The recurrence relation  $a_n = 4a_{n-1} + 5a_{n-2}$  is linear homogeneous

Equation of degree 2.

It associated equation is

$$x^2 = 4x+5$$

Rewriting this as

$$x^2 -4x-5=0$$

$$x^2 -5x+x-5=0$$

$$(x-5)(x+1)=0$$

$$X=5 \text{ or } x=-1$$

The roots of the equation is  $s_1 = 5$  and  $s_2 = -1$

Now, by teorem(i)

We can find value of u and v

$$\text{From } a_n = us_1^n + vs_2^n \quad \text{-----(A)}$$

For n=1

$$a_1 = us_1 + vs_2$$

$$2 = u(5) + v(-1)$$

$$2 = 5u - v \quad \text{-----(i)}$$

For n = 2

$$a_2 = us_1^2 + vs_2^2$$

$$6 = u(5)^2 + v(-1)^2$$

$$6 = 25u + v \quad \text{-----(ii)}$$

Solving equation (i) and (ii)

$$5u - v = 2$$

$$25u + v = 6$$

$$+ \quad + \quad +$$

---

$$30u = 8$$

$$U = 8/30$$

Putting values of u in equation (i)

$$2 = 5(8/30) - v$$

$$2 = (8/6) - v$$

$$2 = 8/6 - 2$$

$$2 = 8 - 12/6$$

$$V = -4/6$$

$$V = -2/3$$

Put value of  $u_1, v_1, s_2$  and  $s_2$  in equation (A)

$$a_n = us_1^n + vs_2^n$$

$$a_n = (8/30)(5)^n + (-2/3)(-1)^n$$

$$a_n = 8/30(5)^n + -2/3(-1)^n$$

∴ Which is required formula.

(a) Define :

(i) Semigroup

(ii) Monoid

(iii) Subsemigroup

(iv) Group Homomorphism.

(i) **Semigroup:-**

Let S be a non-empty set and \* be a binary operation on S. The algebraic system (S, \*) is called a semigroup if the operation \* is

(1) The operation \* is a closed operation on set A.

(2) The operation \* is an associative operation.

Or (S, \*) is a semigroup if for any  $x, y, z \in S$ ,

$$(x * y) * z = x * (y * z)$$

**Free semigroup:**

If  $*$  is an associative binary operation, and  $(A, *)$  is a semigroup. The semigroup  $(A, *)$  is called free semigroup by  $A$ .

Ex:

Consider an algebraic system  $(S, *)$  where  $S = \{1,2,3,5,7,9, \dots\}$  the set of all positive odd integers and  $*$  is a binary operation means multiplication. Determine whether  $(S, *)$  is a semigroup.

**(ii) Monoid:-**

Let us consider an algebraic system  $(M, *)$ , where  $*$  is a binary operation on  $M$ . Then the system  $(M, *)$  is said to be a monoid if it satisfies the following properties:

- (1) The operation  $*$  is a closure operation on set  $A$ .
- (2) The operation  $*$  is an associative operation.
- (3) There exists an identify element w. r. t. The operation  $*$ .

Ex:-

Consider an algebra system  $(N, +)$ , where the set  $N = \{0,1,2,3, \dots\}$  the set of natural numbers and  $+$  is an addition operation. Determine whether  $(N, +)$  is a monoid.

**(iii) Subsemigroup:-**

Let  $(S, *)$  be a semigroup and  $T \subseteq S$ , if the set  $T$  is closed under the operation  $*$  then  $(T, *)$  is said to be subsemigroups of  $(S, *)$ .

Ex:

Consider a semigroup  $(N, +)$ , where  $N$  is the set of all natural number and  $+$  is an addition operation.

The algebraic system  $(E, +)$  is a subsemigroup of  $(N, +)$ , where  $E$  is a set of all +ve even integer.

**(iv) Group homomorphism:-**

Let  $(S, *)$  and  $(T, *)$  be two semigroups. An everywhere defined function  $f: S \rightarrow T$  is called a homomorphism from  $(S, *)$  and  $(T, *)$

$$f(a * b) = f(a) * f(b)$$

For all  $a$  and  $b$  in  $S$ .

If  $f$  is also onto.

We say that  $T$  is a homomorphic image of  $S$ .

(b) Let the number of edges of graph  $G$  be  $m$ . Then  $G$  has a Hamiltonian circuit if  $m \geq \frac{1}{2}(n^2 - 3n + 6)$  where  $n = n_0$  of vertices.

Sol== partial order set :

Let  $A$  is a relation on set  $A$ . then relation  $R$  is called partial order. If it is reflexive, antisymmetric and transitive.



If R is a partial order relation on set A. then set A together with partial order relation R is know as partial orderd set or partial order set.

Ex. Let Z be a set f integers " $\leq$ " be a relation on Z.

∴ Reflexive property is satisfied .

$$(\cdot \cdot a \leq a \quad \forall a \in Z)$$

Let  $a, b \in Z$

$$a \leq b \text{ and } b \leq a \Rightarrow a=b$$

∴ Antisymmetric property is satisfied

$$a \leq b \text{ and } b \leq c \Rightarrow a \leq c$$

∴ Transitive property is satiesfide.

∴ " $\leq$ " is a partial order relation on Z

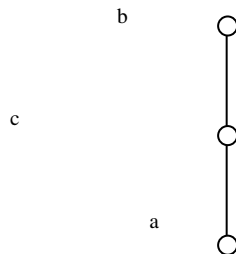
Similarly " $\geq$ " is also a partial order relation on Z.

Chain order set:

If every pair of element in a poset is comparable than poset A is called linear order set .Or set A is chain .

Ex.  $A = \{a, b, c\}$

$$a \leq c, c \leq b$$



This order is in linear or chain.

Hence it is called chain or linear orderd.

Lexicographic:

Let  $A \times B$  is a cartesian product of two sets A & B .we define " $<$ " as follow.

$$(a \ b) < (a' \ b') \text{ if } a < a' \text{ or if } a = a' \text{ then } b < b'$$

This is used in dictionary.

Hence it is also as dictionary

Ex. Help, help

$$\text{Help} < \text{help}$$

Isomorphism:

Let  $(A \leq)$  and  $(A' \leq')$  be posets and let  $f: A \rightarrow A'$  be a one-to-one correspondence between  $f: A \rightarrow A'$  The function f is called an Isomorphism from A to A'

It for any  $a, b \in A, a \leq b$

$$\Leftrightarrow f(a) \leq' a(b).$$

(Proof left)

(c) Let  $G$  be the set of all non-zero real numbers and let  
 $a * b = \frac{ab}{2}$ . Show that  $(G, *)$  is an Abelian  
group.

**Sol== To show:**  $(G, *)$  is an abelian group.

**Closure property:**

The set  $G$  is closed under the operation  $*$ .

Since,  $a * b = \frac{ab}{2}$  is a real number.

Hence, belongs to  $G$ .

**Associative property:**

The operation  $*$  is associative.

Let  $a, b, c \in G$ , then

We have

$$\begin{aligned}(a * b) * c &= \left( \frac{ab}{2} \right) * c \\ &= \frac{(ab)c}{4}\end{aligned}$$

$$= \frac{abc}{4}$$

Similarly,  $a * (b * c) * a = \left( \frac{ab}{2} \right)$

$$= \frac{a(bc)}{4}$$

$$= \frac{abc}{4}$$

(d) Let  $T$  be the set of all even integers. Show that the semigroup  $(\mathbb{Z}, +)$  and  $(T, +)$  are isomorphic

Sol== solu:

Let  $a$  and  $d$  be any element in  $G$ , since  $R$  is an equivalence relation  $b \in [a]$

If and only if  $[b] = [a]$

Also  $G/R$  is a group

Therefore  $[b] = [a]$  if and only if

$$[e] = [a^{-1}] [b]$$

$$= [a^{-1} b]$$

Thus,  $b \in [a]$  if and only if

$$H = [e] = [a^{-1} b]$$

That is,  $b \in [a]$  if and only if

$$a^{-1}b \in H \text{ or } b \in aH$$

This prove that

$$[a] = aH \text{ for every } a \in G$$

We can show

Similar that  $b \in [a]$  if and only if

$$H = [e]$$

$$= [b] [a]^{-1}$$

$$= [ba]^{-1}$$

This is equivalent to the statement  $[a] = Ha$

Thus,  $[a]$  = are isomorphic

Either

1.(A) (a) Let A, B and C be finite sets with  $|A| = 6$ ,  $|B| = 8$ ,  $|C| = 6$ ,  $|A * B * C| = 11$ ,  $|A \cap B| = 3$ ,  $|A \cap C| = 2$  and  $|B \cap C| = 5$ . Find  $|A \cap B \cap C|$ .

Basic of induction: For  $n=1$

$$P(1) = \text{LHS} = A_1 \cap B$$

$$P(1) = \text{RHS} = A_1 \cap B$$

$\therefore \text{LHS} = \text{RHS}$

$$A_1 \cap B = A_1 \cap B$$

P(1) is true for  $n=1$

Induction step: For  $n=k$

$$P(k) = \text{LHS} = \left( \bigcup_{i=1}^k A_i \right) \cap B = \bigcup_{i=1}^k (A_i \cap B)$$

$$P(k) = \text{RHS} = \bigcup_{i=1}^k (A_i \cap B)$$

$\therefore p(k)$  is also true for  $n=k$

Similarly for  $n = k+1$

$$\text{LHS} = \left( \bigcup_{i=1}^{k+1} A_i \right) \cap B$$

$$= (A_1 \cup A_2 \cup \dots \cup A_k \cup A_{k+1}) \cap B$$

$$= \left( \left( \bigcup_{i=1}^k A_i \right) \cup A_{k+1} \right) \cap B$$

$$= \left( \left( \bigcup_{i=1}^k A_i \right) \cap B \right) \cup (A_{k+1} \cap B) \quad \{\text{by distributive property}\}$$

$$= \left( \bigcup_{i=1}^k (A_i \cap B) \right) \cup (A_{k+1} \cap B)$$

(c) Prove by mathematical induction :

3

$$1^2 + 3^2 + 5^2 + \dots + (2n - 1)^2 = n(2n + 1)(2n - 1).$$

$$\text{Sol} \Rightarrow a_n = 4a_{n-1} + 5a_{n-2} \text{ where } a_1 = 2, a_2 = 6$$

Soln: First find sequence for recurrence relation

$$a_n = 4a_{n-1} + 5a_{n-2}$$

$$\text{For } n = 3 \quad a_3 = 4a_{3-1} + 5a_{3-2}$$

$$= 4a_2 + 5a_1$$

$$= 4(6) + 5(2)$$

$$= 24 + 10$$

$$= 34$$

$$\text{For } n = 4 \quad a_4 = 4a_{4-1} + 5a_{4-2}$$

$$= 4a_3 + 5a_2$$

$$= 4(34) + 5(6)$$

$$= 166$$

$$\text{For } n = 5 \quad a_5 = 4a_{5-1} + 5a_{5-2}$$

$$= 4a_4 + 5a_3$$

$$= 4(166) + 5(34)$$

$$= 834$$

$\therefore$  Sequence is 2,6,34,166,834----

The recurrence relation  $a_n = 4a_{n-1} + 5a_{n-2}$  is linear homogeneous

Equation of degree 2.

It associated equation is

$$x^2 = 4x + 5$$

Rewriting this as

$$x^2 - 4x - 5 = 0$$

$$x^2 - 5x + x - 5 = 0$$

$$(x-5)(x+1)=0$$

$$X=5 \text{ or } x=-1$$

The roots of the equation is  $s_1 = 5$  and  $s_2 = -1$

Now, by theorem(i)

We can find value of u and v

$$\text{From } a_n = us_1^n + vs_2^n \text{----- (A)}$$

For  $n=1$

$$a_1 = us_1 + vs_2$$

$$2 = u(5) + v(-1)$$

$$2 = 5u - v \quad \text{-----(i)}$$

For  $n = 2$

$$a_2 = us_1^2 + vs_2^2$$

$$6 = u(5)^2 + v(-1)^2$$

$$6 = 25u + v \quad \text{-----(ii)}$$

Hence proved

(c) Prove that :

If A, B and C are Boolean matrices of compatible sizes then,

$$(A \vee B) \vee C = A \vee (B \vee C)$$

**Sol== Solution:**

Let us assume that

$$A = [a_{ij}] m \times n$$

$$B = [b_{jk}] m \times n$$

$$C = [c_{kl}] m \times n$$

$$(A \vee B) \vee C = A \vee (B \vee C)$$

$$\text{i.e } (A+B)+C = A+(B+C)$$

Now,

$$A + B = [a_{ij}] m \times n + [b_{jk}] m \times n$$

$$= [a_{ij} + b_{jk}] m \times n$$

$$(A + B) + C = [a_{ij} + b_{jk}] m \times n + [c_{kl}] m \times n$$

$$= [a_{ij} + b_{jk} + c_{kl}] m \times n \longrightarrow (1)$$

$$(B + C) = [b_{jk}] m \times n + [c_{kl}] m \times n$$

$$= [b_{jk} + c_{kl}] m \times n$$

$$(B + C) + A = [a_{ij}] m \times n + [b_{jk} + c_{kl}] m \times n$$

$$= [a_{ij} + b_{jk} + c_{kl}] m \times n \longrightarrow (2)$$

From equation (1) & (2)

We get  $(A + B) + C = A + (B + C)$

i.e.

$(A + B) + C = A + (B + C)$
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**EITHER**

2. (a) Obtain principle disjunctive normal form of :

(i)  $(P \vee Q) \wedge (Q \vee R) \wedge (P \vee R)$

(ii)  $P \vee Q \vee R$

$$\Rightarrow [(P \vee Q) \wedge (P \vee R) \wedge (Q \vee R)] \vee \neg[(P \vee R) \wedge (Q \vee R)] \quad \{\text{by associative property}\}$$

$$\Rightarrow [(P \vee Q) \wedge (P \vee R) \wedge (Q \vee R)] \vee [\neg(P \vee R) \vee \neg(Q \vee R)] \quad \{\text{By demorgans and distributive prop. Respt.}\}$$

$$\Rightarrow [(P \vee Q) \wedge (P \vee R) \wedge (Q \vee R)] \vee [\neg(P \vee R) \vee \neg(Q \vee R)] \quad \{\text{by demorgans property}\}$$

$$\Rightarrow [(P \vee Q) \wedge (P \vee R) \wedge (Q \vee R)] \vee \neg(P \vee R) \vee \neg(Q \vee R) \quad \{\text{by demorgans property \& } \neg\neg A \Rightarrow A\}$$

$$\Rightarrow [P \vee (Q \wedge (Q \vee R))] \vee \neg(P \vee R) \vee \neg(Q \vee R) \quad \{\text{by distributive property.}\}$$

$$\Rightarrow [P \vee (Q \wedge Q) \wedge R] \vee \neg(P \vee R) \vee \neg(Q \vee R) \quad \{\text{by associative property}\}$$

$$\Rightarrow (P \vee (Q \wedge R)) \vee \neg(P \vee R) \vee \neg(Q \vee R) \quad \{\text{by idempotent property \& } P \wedge P \Rightarrow P\}$$

$$\Rightarrow T \quad \{P \vee \neg P \Rightarrow T\}$$

Hence Proved



A. Obtain the Principal Disjunction Normal Form of :

1.  $\neg P \vee Q$

2.  $(P \wedge Q) \vee (\neg P \wedge R) \vee (Q \wedge R)$

**Solution:**

1.  $\neg P \vee Q$

$$\Rightarrow (\neg P \wedge T) \vee (Q \wedge T) \quad \{\text{by } P \wedge T = P\}$$

$$\Rightarrow [\neg P \wedge (Q \vee \neg Q)] \vee [Q \wedge (P \vee \neg P)] \quad \{\text{by } P \vee \neg P = T\}$$

$$\Rightarrow [(\neg P \wedge Q) \vee (\neg P \wedge \neg Q)] \vee [(Q \wedge P) \vee (Q \wedge \neg P)] \quad \{\text{by distributive property}\}$$

$$\Rightarrow (\neg P \wedge Q) \vee (\neg P \wedge \neg Q) \vee (Q \wedge P) \vee (Q \wedge \neg P) \quad \{\text{by Associative property}\}$$

$$\Rightarrow (\neg P \wedge Q) \vee (\neg P \wedge \neg Q) \vee (P \wedge Q) \vee (P \wedge \neg Q) \quad \{\text{by Commutative property}\}$$

∴ It is form of sum of elementary product of min term.

Hence, it is in the form of Principal Disjunction Normal Form.

(d) Obtain conjunctive normal form of :

$$\neg(P \vee Q) \quad \Leftrightarrow \quad (P \wedge Q).$$

**Proof:-**



$$\begin{aligned}
& \neg(P \vee Q) \leftrightarrow (P \wedge Q) \\
& \text{by } R \leftrightarrow S \Leftrightarrow (R \rightarrow S) \wedge (S \rightarrow R) \\
& \Leftrightarrow [\neg(P \vee Q) \rightarrow (P \wedge Q)] \wedge [(P \wedge Q) \rightarrow \neg(P \vee Q)] \\
& \Leftrightarrow [\neg\neg(P \vee Q) \vee (P \wedge Q)] \wedge [\neg(P \wedge Q) \vee \neg(P \vee Q)] \quad \{P \rightarrow Q \Rightarrow \neg P \vee Q\} \\
& \Leftrightarrow [(P \vee Q) \vee (P \wedge Q)] \wedge [(\neg P \vee \neg Q) \vee \neg(P \vee Q)] \quad \{\text{By Demorgans property} \& \neg\neg P \Rightarrow P\} \\
& \Leftrightarrow [(P \vee Q \vee P) \wedge (P \vee Q \vee Q)] \wedge [(\neg P \vee \neg Q \vee \neg P) \wedge (\neg P \vee \neg Q \vee \neg Q)] \{\text{By Distributive property}\} \\
& \Leftrightarrow ((P \vee P) \vee Q) \wedge ((Q \vee Q) \vee P) \wedge ((\neg P \vee \neg P) \vee \neg Q) \wedge ((\neg Q \vee \neg Q) \vee \neg P) \{\text{By Associative property}\} \\
& \Leftrightarrow (P \vee Q) \wedge (Q \vee P) \wedge (\neg P \vee \neg Q) \wedge (\neg Q \vee \neg P) \quad \{P \vee P = P\} \\
& \Leftrightarrow (P \vee Q) \wedge (P \vee Q) \wedge (\neg P \vee \neg Q) \wedge (\neg P \vee \neg Q) \quad \{\text{By commulative property}\} \\
& \Leftrightarrow (P \vee Q) \wedge (\neg P \vee \neg Q) \quad \{P \wedge P = P\}
\end{aligned}$$

It is the form of product of elementary sum of min terms.

Hence it is form of Principal Conjunction Normal Form.

(c) Obtain Principal Disjunctive Normal Form of  $P \rightarrow ((P \rightarrow Q) \wedge \neg(\neg Q \vee \neg P))$ .

$$\begin{aligned}
\text{Sol} & \Rightarrow (\neg P \wedge T) \vee (Q \wedge T) && \{\text{by } P \wedge T = P\} \\
& \Rightarrow [\neg P \wedge (Q \vee \neg Q)] \vee [Q \wedge (P \vee \neg P)] && \{\text{by } P \vee \neg P = T\} \\
& \Rightarrow [(\neg P \wedge Q) \vee (\neg P \wedge \neg Q)] \vee [(Q \wedge P) \vee (Q \wedge \neg P)] && \{\text{by distributive property}\} \\
& \Rightarrow (\neg P \wedge Q) \vee (\neg P \wedge \neg Q) \vee (Q \wedge P) \vee (Q \wedge \neg P) && \{\text{by Associative property}\} \\
& \Rightarrow (\neg P \wedge Q) \vee (\neg P \wedge \neg Q) \vee (P \wedge Q) \vee (P \wedge \neg Q) && \{\text{by Commutative property}\}
\end{aligned}$$

∴ It is form of sum of elementary product of min term.

Hence, it is in the form of Principal Disjunction Normal Form.

$$\begin{aligned}
& \neg(P \vee Q) \leftrightarrow (P \wedge Q) \\
& \text{by } R \leftrightarrow S \Leftrightarrow (R \rightarrow S) \wedge (S \rightarrow R) \\
& \Leftrightarrow [\neg(P \vee Q) \rightarrow (P \wedge Q)] \wedge [(P \wedge Q) \rightarrow \neg(P \vee Q)] \\
& \Leftrightarrow [\neg\neg(P \vee Q) \vee (P \wedge Q)] \wedge [\neg(P \wedge Q) \vee \neg(P \vee Q)] \quad \{P \rightarrow Q \Rightarrow \neg P \vee Q\} \\
& \Leftrightarrow [(P \vee Q) \vee (P \wedge Q)] \wedge [(\neg P \vee \neg Q) \vee \neg(P \vee Q)] \quad \{\text{By Demorgans property} \& \neg\neg P \Rightarrow P\} \\
& \Leftrightarrow [(P \vee Q \vee P) \wedge (P \vee Q \vee Q)] \wedge [(\neg P \vee \neg Q \vee \neg P) \wedge (\neg P \vee \neg Q \vee \neg Q)] \{\text{By Distributive property}\} \\
& \Leftrightarrow ((P \vee P) \vee Q) \wedge ((Q \vee Q) \vee P) \wedge ((\neg P \vee \neg P) \vee \neg Q) \wedge ((\neg Q \vee \neg Q) \vee \neg P) \{\text{By Associative property}\} \\
& \Leftrightarrow (P \vee Q) \wedge (Q \vee P) \wedge (\neg P \vee \neg Q) \wedge (\neg Q \vee \neg P) \quad \{P \vee P = P\} \\
& \Leftrightarrow (P \vee Q) \wedge (P \vee Q) \wedge (\neg P \vee \neg Q) \wedge (\neg P \vee \neg Q) \quad \{\text{By commulative property}\} \\
& \Leftrightarrow (P \vee Q) \wedge (\neg P \vee \neg Q) \quad \{P \wedge P = P\}
\end{aligned}$$

(b) Find an explicit formula for the sequence defined by  $C_n = 6C_{n-1} + 7C_{n-2}$  with initial conditions

$$C_0 = 2, C_1 = 1.$$

Sol== First find sequence for recurrence relation

$$a_n = 4a_{n-1} + 5a_{n-2}$$

$$\text{For } n = 3 \quad a_3 = 4a_{3-1} + 5a_{3-2}$$

$$= 4a_2 + 5a_1$$

$$= 4(6) + 5(2)$$

$$= 24 + 10$$

$$= 34$$

$$\text{For } n = 4 \quad a_4 = 4a_{4-1} + 5a_{4-2}$$

$$= 4a_3 + 5a_2$$

$$= 4(34) + 5(6)$$

$$= 166$$

$$\text{For } n = 5 \quad a_5 = 4a_{5-1} + 5a_{5-2}$$

$$= 4a_4 + 5a_3$$

$$= 4(166) + 5(34)$$

$$= 834$$

$\therefore$  Sequence is 2,6,34,166,834----

The recurrence relation  $a_n = 4a_{n-1} + 5a_{n-2}$  is linear homogeneous

Equation of degree 2.

It associated equation is

$$x^2 = 4x + 5$$

Rewriting this as

$$x^2 - 4x - 5 = 0$$

$$x^2 - 5x + x - 5 = 0$$

$$(x-5)(x+1) = 0$$

$$X = 5 \text{ or } x = -1$$

The roots of the equation is  $s_1 = 5$  and  $s_2 = -1$

Now, by theorem(i)

We can find value of u and v

$$\text{From } a_n = us_1^n + vs_2^n \quad \text{-----(A)}$$

For  $n=1$

$$a_1 = us_1 + vs_2$$

$$2 = u(5) + v(-1)$$

$$2 = 5u - v \quad \text{-----(i)}$$

For  $n = 2$

$$a_2 = us_1^2 + vs_2^2$$

$$6 = u(5)^2 + v(-1)^2$$

$$6 = 25u + v \quad \text{-----(ii)}$$

Solving equation (i) and (ii)

$$5u - v = 2$$

$$25u + v = 6$$

$$+ \quad + \quad +$$

---

$$30u = 8$$

$$\boxed{U=8/30}$$

Putting values of u in equation (i)

$$2 = 5(8/30) - v$$

$$2 = (8/6) - v$$

$$2 = 8/6 - 2$$

$$2 = 8 - 12/6$$

$$V = -4/6$$

$$\boxed{V = -2/3}$$

Put value of  $u_1, v_1, s_2$  and  $s_2$  in equation (A)

$$a_n = us_1^n + vs_2^n$$

$$a_n = (8/30) (5)^n + (-2/3) (-1)^n$$

$$a_n = 8/30 (5)^n + -2/3 (-1)^n$$

∴ Which is required formula.

(a) Define :

(i) Semigroup

(ii) Monoid

(iii) Subsemigroup

(iv) Group Homomorphism.

**(i) Semigroup:-**

Let S be a non-empty set and \* be a binary operation on S. The algebraic system (S, \*) is called a semigroup if the operation \* is

(1) The operation \* is a closed operation on set A.

(2) The operation \* is an associative operation.

Or (S, \*) is a semigroup if for any x, y, z ∈ S,

$$(x * y) * z = x * (y * z)$$

**Free semigroup:**

If \* is an associative binary operation, and (A, \*) is a semigroup. The semigroup(A, \*) is called free semigroup by A.

Ex:

Consider an algebraic system (S, \*) where S = {1,2,3,5,7,9----} the set of all positive odd integers and \* is a binary operation means multiplication. Determine whether (S, \*) is a semigroup.

**(ii) Monoid:-**

Let us consider an algebraic system (M, \*), where \* is a binary operation on M. Then the system (M, \*) is said to be a monoid if it satisfies the following properties:

(1) The operation \* is a closure operation on set A.

(2) The operation \* is an associative operation.

(3) There exists an identify element w. r. t. The operation \*.

Ex:-

Consider an algebra system (N, +), where the set N = {0,1,2,3----} the set of natural numbers and + is an addition operation. Determine whether (N,+) is a monoid.

**(iii) Subsemigroup:-**

Let  $(S, *)$  be a semigroup and  $T \subseteq S$ , if the set  $T$  is closed under the operation  $*$  then  $(T, *)$  is said to be subsemigroups of  $(S, *)$ .

**Ex:**

Consider a semigroup  $(N, +)$ , where  $N$  is the set of all natural number and  $+$  is an addition operation.

The algebraic system  $(E, +)$  is a subsemigroup of  $(N, +)$ , where  $E$  is a set of all +ve even integer.

**(iv) Group homomorphism:-**

Let  $(S, *)$  and  $(T, *)$  be two semigroups. An everywhere defined function  $f: S \rightarrow T$  is called a homomorphism from  $(S, *)$  and  $(T, *)$

$$f(a * b) = f(a) * f(b)$$

For all  $a$  and  $b$  in  $S$ .

If  $f$  is also onto.

We say that  $T$  is a homomorphic image of  $S$ .

(b) Let the number of edges of graph  $G$  be  $m$ . Then  $G$  has a Hamiltonian circuit if  $m \geq 1/2 (n^2 - 3n + 6)$  where  $n = n_0$  of vertices.

Sol== partial order set :

Let  $A$  is a relation on set  $A$ . then relation  $R$  is called partial order. If it is reflexive, antisymmetric and transitive.

If  $R$  is a partial order relation on set  $A$ . then set  $A$  together with partial order relation  $R$  is known as partial ordered set or partial order set.

Ex. Let  $Z$  be a set of integers " $\leq$ " be a relation on  $Z$ .

∴ Reflexive property is satisfied .

$$(\cdot \cdot a \leq a \quad \forall a \in Z)$$

Let  $a, b \in Z$

$$a \leq b \text{ and } b \leq a \Rightarrow a = b$$

∴ Antisymmetric property is satisfied

$$a \leq b \text{ and } b \leq c \Rightarrow a \leq c$$

∴ Transitive property is satisfied.

∴ " $\leq$ " is a partial order relation on  $Z$

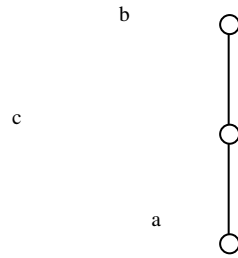
Similarly " $\geq$ " is also a partial order relation on  $Z$ .

Chain order set:

If every pair of element in a poset is comparable then poset  $A$  is called linear order set .Or set  $A$  is chain .

Ex.  $A = \{a, b, c\}$

$$a \leq c, c \leq b$$



This order is in linear or chain.

Hence it is called chain or linear order.

#### Lexicographic:

Let  $A \times B$  is a cartesian product of two sets  $A$  &  $B$ . we define " $<$ " as follow.

$$(a, b) < (a', b') \text{ if } a < a' \text{ or if } a = a' \text{ then } b < b'$$

This is used in dictionary.

Hence it is also as dictionary

Ex. Help, help

Help < help

#### Isomorphism:

Let  $(A, \leq)$  and  $(A', \leq')$  be posets and let  $f: A \rightarrow A'$  be a one-to-one correspondence between  $f: A \rightarrow A'$ . The function  $f$  is called an Isomorphism from  $A$  to  $A'$ .

It for any  $a, b \in A, a \leq b$

$$\Leftrightarrow f(a) \leq' f(b).$$

**(Proof left)**

(c) Let  $G$  be the set of all non-zero real numbers and let

$a * b = ab$ . Show that  $(G, *)$  is an Abelian group.

**Sol== To show:**  $(G, *)$  is an abelian group.

#### **Closure property:**

The set  $G$  is closed under the operation  $*$ .

Since,  $a*b = \frac{ab}{2}$  is a real number.

Hence, belongs to G.

**Associative property:**

The operation \* is associative.

Let a, b, c  $\in$  G, then

We have

$$\begin{aligned}(a*b)*c &= \left(\frac{ab}{2}\right)*c \\ &= \frac{(ab)c}{4}\end{aligned}$$

$$= \frac{abc}{4}$$

$$\begin{aligned}\text{Similarly, } a*(b*c)*a &= \left(\frac{ab}{2}\right) \\ &= \frac{a(bc)}{4} \\ &= \frac{abc}{4}\end{aligned}$$

(d) Let T be the set of all even integers. Show that the semigroup (Z, +) and (T, +) are isomorphic

Sol== solu:

Let a and d be any element in G, since R is an equivalence relation  $b \in [a]$

If and only if  $[b] = [a]$

Also G/R is a group



Therefore  $[b] = [a]$  if and only if

$$[e] = [a^{-1}] [b]$$

$$= [a^{-1} b]$$

Thus,  $b \in [a]$  if and only if

$$H = [e] = [a^{-1} b]$$

That is,  $b \in [a]$  if and only if

$$a^{-1}b \in H \text{ or } b \in aH$$

This prove that

$$[a] = aH \text{ for every } a \in G$$

We can show

Similar that  $b \in [a]$  if and only if

$$H = [e]$$

$$= [b] [a]^{-1}$$

$$= [ba]^{-1}$$

This is equivalent to the statement  $[a] = Ha$

Thus,  $[a]$  = are isomorphic

1] (A) Let  $U = \{a, b, c, d, e, f, g, h, k\}$ ,  $A = \{a, b, c, g\}$ ,  $B = \{d, e, f, g\}$ ,  $C = \{a, c, f\}$  and  $D = \{f, h, k\}$ .

Compute:

- (i)  $A \oplus B$       (ii)  $C \oplus D$   
 (iii)  $\overline{A \cup B}$     (iv)  $\overline{C \cap D}$

Solution:

(B) (i) Construct truth table for statement:

$$p \Rightarrow q \Leftrightarrow \neg p \vee q$$

soln :

(1)    (2) (3)      (4) (5)

P	Q	$(P \Rightarrow Q)$	$\neg P$	$\neg P \vee Q$
T	T	T	F	T
T	F	F	F	F
F	T	T	T	T
F	F	T	T	T

From (3) and (5) column

$\therefore p \Rightarrow q \Leftrightarrow \neg p \vee q$  Hence proved.

(ii) Prove that if  $m^2$  is odd then  $m$  is odd.

Soln:

OR

(c) Explain principle of mathematical induction and use induction method to

Prove:

$$\left( \bigcup_{i=1}^n A_i \right) \cap B = \bigcup_{i=1}^n (A_i \cap B)$$

Soln:

We can prove by mathematical induction.

Basic of induction: For  $n=1$

$$P(1) = \text{LHS} = A_1 \cap B$$

$$P(1) = \text{RHS} = A_1 \cap B$$

$\therefore \text{LHS} = \text{RHS}$

$$A_1 \cap B = A_1 \cap B$$

$P(1)$  is true for  $n=1$

Induction step: For  $n=k$

$$P(k) = \text{LHS} = \left( \bigcup_{i=1}^k A_i \right) \cap B = \bigcup_{i=1}^k (A_i \cap B)$$

$$P(k) = \text{RHS} = \bigcup_{i=1}^k (A_i \cap B)$$

$\therefore P(k)$  is also true for  $n=k$

Similarly for  $n = k+1$

$$\begin{aligned} \text{LHS} &= \left( \bigcup_{i=1}^{k+1} A_i \right) \cap B \\ &= (A_1 \cup A_2 \cup \dots \cup A_k \cup A_{k+1}) \cap B \\ &= \left( \left( \bigcup_{i=1}^k A_i \right) \cup A_{k+1} \right) \cap B \\ &= \left( \left( \bigcup_{i=1}^k A_i \right) \cap B \right) \cup (A_{k+1} \cap B) \quad \{\text{by distributive property}\} \\ &= \left( \bigcup_{i=1}^k (A_i \cap B) \right) \cup (A_{k+1} \cap B) \\ &= \bigcup_{i=1}^{k+1} (A_i \cap B) \\ &= \text{RHS} \end{aligned}$$

Thus the implication  $P(k) \rightarrow P(k+1)$  is a tautology

$\therefore$  By principle of Mathematic Induction

$P(n)$  is true for all  $n \geq 1$

Hence proved

$$\boxed{\left( \bigcup_{i=1}^n A_i \right) \cap B = \bigcup_{i=1}^n (A_i \cap B)}$$

(D) Solve the recurrence relation:

$$a_n = 4a_{n-1} + 5a_{n-2} \quad \text{where } a_1=2, a_2=6$$

Soln: First find sequence for recurrence relation

$$a_n = 4a_{n-1} + 5a_{n-2}$$

$$\text{For } n=3 \quad a_3 = 4a_{3-1} + 5a_{3-2}$$

$$\begin{aligned} &= 4a_2 + 5a_1 \\ &= 4(6) + 5(2) \\ &= 24 + 10 \\ &= 34 \end{aligned}$$

$$\text{For } n=4 \quad a_4 = 4a_{4-1} + 5a_{4-2}$$

$$\begin{aligned} &= 4a_3 + 5a_2 \\ &= 4(34) + 5(6) \\ &= 166 \end{aligned}$$

$$\text{For } n=5 \quad a_5 = 4a_{5-1} + 5a_{5-2}$$

$$\begin{aligned} &= 4a_4 + 5a_3 \\ &= 4(166) + 5(34) \\ &= 834 \end{aligned}$$

∴ Sequence is 2,6,34,166,834----

The recurrence relation  $a_n = 4a_{n-1} + 5a_{n-2}$  is linear homogeneous

Equation of degree 2.

It associated equation is

$$x^2 = 4x+5$$

Rewriting this as

$$x^2 -4x-5=0$$

$$x^2 -5x+x-5=0$$

$$(x-5)(x+1)=0$$

$$X=5 \text{ or } x=-1$$

The roots of the equation is  $s_1 =5$  and  $s_2 =-1$

Now, by theorem(i)

We can find value of u and v

$$\text{From } a_n = us_1^n + vs_2^n \text{----- (A)}$$

For n=1

$$a_1 = us_1 + vs_2$$

$$2 = u(5) + v(-1)$$

$$2 = 5u - v \quad \text{-----(i)}$$

For n = 2

$$a_2 = us_1^2 + vs_2^2$$

$$6 = u(5)^2 + v(-1)^2$$

$$6 = 25u + v \quad \text{-----(ii)}$$

Solving equation (i) and (ii)

$$5u - v = 2$$

$$25u + v = 6$$

$$+ \quad + \quad +$$

---

$$30u = 8$$

$U=8/30$
----------

Putting values of u in equation (i)

$$2 = 5(8/30) - v$$

$$2 = (8/6) - v$$

$$2 = 8/6 - 2$$

$$2 = 8 - 12/6$$

$$V = -4/6$$

$V = -2/3$
------------

Luke of

Put value of  $u_1, v_1, s_2$  and  $s_2$  in equation (A)

$$a_n = us_1^n + vs_2^n$$

$$a_n = (8/30)(5)^n + (-2/3)(-1)^n$$

$$a_n = 8/30(5)^n + -2/3(-1)^n$$

∴ Which is required formula?

## Q.2

Either

**A. Explain dual formula and show that if  $A(P,Q,R) = \neg P \wedge \neg(Q \vee R)$  then**

- i.  $A(\neg P, \neg Q, \neg R) \Leftrightarrow \neg A^*(P, Q, R)$**
- ii.  $\neg A(P, Q, R) \Leftrightarrow A^*(\neg P, \neg Q, \neg R)$**

**Solution:**

$$(i) A(\neg P, \neg Q, \neg R) \Leftrightarrow \neg A^*(P, Q, R)$$

We have to prove that, if  $A(P,Q,R) = \neg P \wedge \neg(Q \vee R)$  then

$$A(\neg P, \neg Q, \neg R) \Leftrightarrow \neg A^*(P, Q, R)$$

$$\Rightarrow A(P, Q, R) = \neg P \wedge \neg(Q \vee R)$$

$$= (\neg P \wedge \neg Q \wedge \neg R)$$

$$= \neg(P \vee Q \vee R)$$

$$\Rightarrow A(\neg P, \neg Q, \neg R) = \neg(P \vee Q \vee R) \quad \xrightarrow{(1)}$$

$$\Rightarrow A^*(P, Q, R) = \neg P \wedge \neg(Q \vee R)$$

$$= (\neg P \wedge \neg Q \wedge \neg R)$$

$$= \neg(P \vee Q \vee R)$$

$$\Rightarrow \neg A^*(P, Q, R) = \neg(\neg(P \vee Q \vee R)) \quad \xrightarrow{(2)}$$

From eq 1 & 2

We get,

$A(\neg P, \neg Q, \neg R) \Leftrightarrow \neg A^*(P, Q, R)$
---

Hence proved

(ii) We have to prove that, if  $A(P,Q,R) = \neg P \wedge \neg(Q \vee R)$  then

$$\neg A(P, Q, R) \Leftrightarrow A^*(\neg P, \neg Q, \neg R)$$

$$\begin{aligned}
&\Rightarrow A(P,Q,R) = \neg P \wedge \neg(Q \vee R) \\
&= \neg(P \vee (Q \vee R)) \\
\Rightarrow \neg A(P,Q,R) &= \neg\neg(P \vee (Q \vee R)) \\
&= P \vee (Q \vee R) \\
\Rightarrow A^*(P,Q,R) &= \neg P \vee \neg(Q \wedge R) \quad \longrightarrow (1) \\
\Rightarrow A^*(\neg P, \neg Q, \neg R) &= P \vee \neg(\neg Q \wedge \neg R) \\
&= P \vee \neg(\neg(Q \wedge R)) \\
\Rightarrow A^*(\neg P, \neg Q, \neg R) &= P \vee (Q \vee R) \quad \longrightarrow (2)
\end{aligned}$$

From eq 1 & 2

We get,

$$\neg A(P,Q,R) \quad \boxed{\neg A(P,Q,R) \leftrightarrow A^*(\neg P, \neg Q, \neg R)}$$

Hence proved

**A. Obtain the Principal Disjunction Normal Form of :**

1.  $\neg P \vee Q$
2.  $(P \wedge Q) \vee (\neg P \wedge R) \vee (Q \wedge R)$

**Solution:**

$$\begin{aligned}
&1. \quad \neg P \vee Q \\
\Rightarrow &(\neg P \wedge T) \vee (Q \wedge T) && \{\text{by } P \wedge T = P\} \\
\Rightarrow &[\neg P \wedge (Q \vee \neg Q)] \vee [Q \wedge (P \vee \neg P)] && \{\text{by } P \vee \neg P = T\} \\
\Rightarrow &[(\neg P \wedge Q) \vee (\neg P \wedge \neg Q)] \vee [(Q \wedge P) \vee (Q \wedge \neg P)] && \{\text{by distributive property}\} \\
\Rightarrow &(\neg P \wedge Q) \vee (\neg P \wedge \neg Q) \vee (Q \wedge P) \vee (Q \wedge \neg P) && \{\text{by Associative property}\} \\
\Rightarrow &(\neg P \wedge Q) \vee (\neg P \wedge \neg Q) \vee (P \wedge Q) \vee (P \wedge \neg Q) && \{\text{by Commutative property}\}
\end{aligned}$$

∴ It is form of sum of elementary product of min term.

Hence, it is in the form of Principal Disjunction Normal Form.

$$\begin{aligned}
&2. \quad (P \wedge Q) \vee (\neg P \wedge R) \vee (Q \wedge R) \\
\Rightarrow &[(P \wedge Q) \wedge T] \vee [(\neg P \wedge R) \wedge T] \vee [(Q \wedge R) \wedge T] \\
&\{\text{By } P \vee \neg P = T\} \\
\Rightarrow &[(P \wedge Q) \wedge (R \vee \neg R)] \vee [(\neg P \wedge R) \wedge (Q \vee \neg Q)] \vee [(Q \wedge R) \wedge (P \vee \neg P)] && \{\text{by } P \vee \neg P = T\} \\
\Rightarrow &[(P \wedge Q \wedge R) \vee (P \wedge Q \wedge \neg R)] \vee [(\neg P \wedge R \wedge Q) \vee (\neg P \wedge R \wedge \neg Q)] \vee [(Q \wedge R \wedge P) \vee (Q \wedge R \wedge \neg P)] \\
&\text{by distributive property}
\end{aligned}$$

{

$$\Rightarrow (P \wedge Q \wedge R) \vee (P \wedge Q \wedge \neg R) \vee (\neg P \wedge Q \wedge R) \vee (\neg P \wedge \neg Q \wedge R) \vee (P \wedge Q \wedge R) \vee (\neg P \wedge Q \wedge R)$$

{by Commutative property}

$$\Rightarrow (P \wedge Q \wedge R) \vee (P \wedge Q \wedge \neg R) \vee (\neg P \wedge Q \wedge R) \vee (\neg P \wedge \neg Q \wedge R) \quad \{\text{by } P \vee \neg P = P\}$$

$\therefore$  it is form of sum of elementary product of min term.

Hence, It is in the form of Principal Disjunction Normal Form.

**OR**

**(C) Determine whether the conclusion C follows logically from premises H1 and H2:**

- (i) **H1 :  $P \rightarrow Q$    H2 :  $\neg P$    C :  $Q$**
- (ii) **H1:  $P \rightarrow Q$    H2:  $\neg(P \wedge Q)$  C:  $\neg P$**
- (iii) **H1:  $\neg P$    H2:  $P \Leftrightarrow Q$    C:  $\neg(P \wedge Q)$**
- (iv) **H1:  $P \rightarrow Q$    H2:  $Q$    C:  $P$**

**(D) Show that:**

$$(X) (P(X) \vee Q(X)) \Rightarrow (X) P(X) \vee (\exists X) Q(X)$$

**Q.3)**

**(A) (i) If  $A \subseteq C$  and  $B \subseteq D$  then prove that**

$$A \times B \subseteq C \times D$$

Soln :-

Given that If  $A \subseteq C$  and  $B \subseteq D$

To prove:  $A \times B \subseteq C \times D$

Proof: Let  $(x,y) \in A \times B$ , then  $x \in A$  and  $y \in B$

Since  $A \subseteq C$  and  $B \subseteq D$ ,  $x \in C$  and  $y \in D$

Hence,  $(x,y) \in C \times D$

Then

$A \times B \subseteq C \times D$
-----------------------------------

Hence proved

(iii) Let  $A = \{1,2,4\}$ ,  $B = \{2,5,7\}$  and  $C = \{1,3,7\}$

Show that  $A \times (B \cap C) = (A \times B) \cap (A \times C)$

**(B) (ii) Let R and S are relation from A to B then prove that:**

(i)  $(R \cap S)^{-1} = R^{-1} \cap S^{-1}$

**Ans:** (i) Let  $(a,b) \in (R \cap S)^{-1}$

So, we have  $(b,a) \in (R \cap S)$

Now  $(b,a) \in R$  and  $(b,a) \in S$

This mean  $(a,b) \in R^{-1}$  and  $(a,b) \in S^{-1}$

Hence  $(a,b) \in R^{-1} \cap S^{-1}$  -----(i)

Conversely,

Let,  $(a,b) \in R^{-1} \cap S^{-1}$

So, we have  $(a,b) \in R^{-1}$  and  $(a,b) \in S^{-1}$

This means  $(b,a) \in R$  and  $(b,a) \in S$

So,  $(b,a) \in (R \cap S)$  -----(ii)

From (i) and (ii) we have

$$(R \cap S)^{-1} = R^{-1} \cap S^{-1}$$

Hence Proved

$$(ii) \overline{R \cap S} = \overline{R} \cup \overline{S}$$

OR

( C ) Let  $A = B = \{1,2,3,4\}$ ,  $R = \{(1,1),(1,3),(2,3),(3,1),(4,2),(4,4)\}$

And  $S = \{(1,2), (2,3),(3,1), (3,2), (3,1), (3,2), (4,3)\}$  then compute

$$M_{R \cap S}, M_{R \cup S}, {}^M R^{-1}, {}^M S^{-1}.$$

( D ) Let  $A$  be a set with  $|A| = n$  and let  $R$  be a relation on  $A$ . Then prove that:

$$R^\infty = R \cup R^2 \cup \dots \cup R^n$$

EITHER

4 (A) (i) Explain:

Post,chain, Hasses Diagram and draw the Hasses diagram of posset  $A$  with

5 (A) Let  $(G, *)$  and  $(G', *')$  be two group and let  $f : G \rightarrow G'$  be a homomorphism from  $G$  to  $G'$  then

(i)  $f(e) = e'$  where  $e$  is identity of  $G$  and  $e'$  is identity of  $G'$

(ii)  $f(a^{-1}) = (f(a))^{-1}$

solun: (a) Let  $x = f(e)$  then

$$\begin{aligned} x *' x &= f(e) *' f(e) \\ &= f(e * e) \\ &= f(e) \\ &= x \end{aligned}$$

So,  $x *' x = x$

Multiplying both side by  $x^{-1}$

On right, we obtain

$$X = x *' x *' x^{-1} = x *' x^{-1} = e'$$

Thus  $f(e) = e'$



$$(b) a * a^{-1} = e$$

$$\text{So, } f(a * a^{-1}) = f(e)$$

$$= e' \text{ by part(a)}$$

$$\text{Or } f(a) * f(a^{-1}) = e's$$

Since,  $f$  is a homomorphism

$$\text{Similarly, } f(a^{-1}) * f(a) = e'$$

$$\text{Hence } f(a^{-1}) = (f(a))^{-1}$$

**(B) (i) Let  $G$  be an abelian group with identity  $e$  and let  $H = \{x : x^2 = e\}$ . Show that  $H$  is a subgroup of  $G$ .**

**(ii) Insertion of two sub subgroup of  $G$  is a subgroup of  $G$**

**OR**

**(C) Define:**

**Finite state Machine, state transition function and Moore Machine and**

**Draw diagraph whose table is: (summer-13)**

**(D) Let  $R$  be a congruence relation on a group  $G$  and let  $H = [e]$ , the equivalent class containing the identity. Then  $H$  is a normal subgroup of  $G$  and for each**

**$a \in G [a] = Ha = aH$  prove this**

solu:

Let  $a$  and  $d$  be any element in  $G$ , since  $R$  is an equivalence relation  $b \in [a]$

If and only if  $[b] = [a]$

Also  $G/R$  is a group

Therefore  $[b] = [a]$  if and only if

$$[e] = [a^{-1}] [b]$$

$$= [a^{-1} b]$$

Thus,  $b \in [a]$  if and only if

$$H = [e] = [a^{-1} b]$$

That is,  $b \in [a]$  if and only if

$$a^{-1}b \in H \text{ or } b \in aH$$

This prove that

$$[a] = aH \text{ for every } a \in G$$

We can show

Similar that  $b \in [a]$  if and only if

$$H = [e]$$

$$= [b] [a]^{-1}$$

$$= [ba]^{-1}$$

This is equivalent to the statement  $[a] = Ha$

Thus,  $[a] = aH = Ha$  and  $H$  is normal.

## 1. EITHER

(A) Prove by mathematical induction that, for all

$$n! \geq 1,$$

$$n! \geq 2^{n-1}, \text{ where}$$

$$1! = 1 \text{ and } n! = n(n-1)!$$

**Solution:-**

(B) If a and b are two positive integers, then prove that

$$\text{GCD}(a,b) \cdot \text{LCM}(a,b) = ab$$

Verify above result for a = 100, b = 80

**Proof:-**

Let  $p_1, p_2, p_3, \dots, p_n$  be the prime factor of a and b.

$$\left. \begin{aligned} \therefore a &= p_1^{a_1} * p_2^{a_2} * p_3^{a_3} * \dots * p_n^{a_n} \\ \therefore b &= p_1^{b_1} * p_2^{b_2} * p_3^{b_3} * \dots * p_n^{b_n} \end{aligned} \right\} \rightarrow (1)$$

By definition of LCM and GCD, we get

$$\text{GCD}(a, b) = p_1^{\min(a_1, b_1)} * p_2^{\min(a_2, b_2)} * p_3^{\min(a_3, b_3)} * \dots * p_n^{\min(a_n, b_n)}$$

$$\text{LCM}(a, b) = p_1^{\max(a_1, b_1)} * p_2^{\max(a_2, b_2)} * p_3^{\max(a_3, b_3)} * \dots * p_n^{\max(a_n, b_n)}$$

$$L.H.S = \text{GCD}(a, b) * \text{LCM}(a, b)$$

$$= [ p_1^{\min(a_1, b_1)} * p_2^{\min(a_2, b_2)} * p_3^{\min(a_3, b_3)} * \dots * p_n^{\min(a_n, b_n)} ] *$$

$$[ p_1^{\max(a_1, b_1)} * p_2^{\max(a_2, b_2)} * p_3^{\max(a_3, b_3)} * \dots * p_n^{\max(a_n, b_n)} ]$$

$$= [ p_1^{\min(a_1, b_1)} * p_1^{\max(a_1, b_1)} ] * [ p_2^{\min(a_2, b_2)} * p_2^{\max(a_2, b_2)} ] *$$

$$[ p_3^{\min(a_3, b_3)} * p_3^{\max(a_3, b_3)} ] * \dots * [ p_n^{\min(a_n, b_n)} * p_n^{\max(a_n, b_n)} ]$$

$$= (p_1^{a_1} * p_1^{b_1}) * (p_2^{a_2} * p_2^{b_2}) * (p_3^{a_3} * p_3^{b_3}) * \dots * (p_n^{a_n} * p_n^{b_n})$$

$$= (p_1^{a_1} * p_2^{a_2} * p_3^{a_3} * \dots * p_n^{a_n}) * (p_1^{b_1} * p_2^{b_2} * p_3^{b_3} * \dots * p_n^{b_n})$$

$$\text{form eq}^n (1)$$

$$= a * b$$

$$= R.H.S$$

$$\therefore L.H.S = R.H.S$$

$$\therefore \text{LCM}(a, b) \cdot \text{GCD}(a, b) = a \cdot b$$

**Verification of this result for a = 100, b = 80**

$$100 = 1 * 2 * 2 * 5 * 5$$

$$80 = 1 * 2 * 2 * 2 * 2 * 5$$

$$\text{GCD}(100, 80) = 1 * 2 * 2 * 5 = 20$$

$$= 20$$

$$\text{LCM}(100, 80) = 1 * 2 * 2 * 2 * 2 * 5 * 5$$

$$= 400$$

$$\text{LCM}(100,80) \cdot \text{GCD}(100,80) = 100 \cdot 80$$

$$20 \cdot 400 = 8000$$

$$8000 = 8000$$

**OR**

**(C) Prove that :-**

**If A, B and C are Boolean matrices of compatible sizes, then**

$$(A \oplus B) \ominus C = A \ominus (B \ominus C)$$

**Solution:**

Let us assume that

$$A = [a_{ij}] m \times n$$

$$B = [b_{jk}] m \times n$$

$$C = [c_{kl}] m \times n$$

$$(A \vee B) \vee C = A \vee (B \vee C)$$

$$\text{i.e. } (A+B)+C = A+(B+C)$$

Now,

$$A + B = [a_{ij}] m \times n + [b_{jk}] m \times n$$

$$= [a_{ij} + b_{jk}] m \times n$$

$$(A + B) + C = [a_{ij} + b_{jk}] m \times n + [c_{kl}] m \times n$$

$$= [a_{ij} + b_{jk} + c_{kl}] m \times n \quad \xrightarrow{(1)}$$

$$(B + C) = [b_{jk} + c_{kl}] m \times n$$

$$= [b_{jk} + c_{kl}] m \times n$$

$$(B + C) + A = [a_{ij}] m \times n + [b_{jk} + c_{kl}] m \times n$$

$$= [a_{ij} + b_{jk} + c_{kl}] m \times n \quad \xrightarrow{(2)}$$

From equation (1) & (2)

We get  $(A \vee B) \vee C = A \vee (B \vee C)$

i.e.

$$(A \oplus B) \ominus C = A \ominus (B \ominus C)$$

**(D) Let m and n be integers. Prove that  $n^2 = m^2$**

**If and only if n is m or n is -m.**

**2. EITHER**

(A) Prove the equivalence.

$$(P \vee Q) \wedge (\neg P \wedge Q) \Leftrightarrow (\neg P \wedge Q) \text{ without using truth table.}$$

**Proof:-**  $(P \vee Q) \wedge (\neg P \wedge Q) \Leftrightarrow (\neg P \wedge Q)$

L.H.S

$$\begin{aligned} &\Rightarrow (P \vee Q) \wedge (\neg P \wedge (\neg P \wedge Q)) \\ &\Rightarrow (P \vee Q) \wedge ((\neg P \wedge \neg P) \wedge Q) && \{ \text{By associative properties} \} \\ &\Rightarrow (P \vee Q) \wedge (\neg P \wedge Q) && \{ \text{By idempotent properties} \} \\ &\Rightarrow (P \wedge (\neg P \wedge Q)) \vee (Q \wedge (\neg P \wedge Q)) && \{ \text{By distributive properties} \} \\ &\Rightarrow ((P \wedge \neg P) \wedge Q) \vee ((Q \wedge Q) \wedge \neg P) && \{ \text{By associative properties} \} \\ &\Rightarrow (F \wedge Q) \vee (Q \wedge \neg P) && \{ P \wedge \neg P = F \text{ \& idempotent properties} \} \\ &\Rightarrow F \vee (Q \wedge \neg P) && \{ F \wedge Q = F \} \\ &\Rightarrow F \vee (\neg P \wedge Q) && \{ \text{By commutative properties} \} \\ &\Rightarrow (\neg P \wedge Q) && \{ P \vee F = P \} \\ &\Rightarrow \text{R.H.S} \end{aligned}$$

hence

$$(P \vee Q) \wedge (\neg P \wedge Q) \Leftrightarrow (\neg P \wedge Q)$$

(B) Show that the following premises are inconsistent:-

- (i) If Jack misses many classes through illness, then he fail high school.
- (ii) If Jack fails high school, the he is uneducated.
- (iii) If Jack reads a lot of books, then he is not uneducated.
- (iv) Jack misses many classes through illness and reads a lot of books.

**Proof:-**

We have to prove that given premises are inconsistent. To prove inconsistent we have to derive contradiction from the given premises.

Let, E: Jack misses many classes.

S: Jack fails high School.

A: Jack read a lot of books.

H: Jack is uneducated.

Given premises are,

$$\begin{aligned} E &\rightarrow S \\ S &\rightarrow H \\ A &\rightarrow \neg H \\ E &\rightarrow A \end{aligned}$$

{1}	(1) $E \rightarrow S$	{Rule P}
{2}	(2) $S \rightarrow H$	{Rule P}
{1,2}	(3) $E \rightarrow H$	{Rule T i, e $P \rightarrow Q, Q \rightarrow R \Rightarrow P \rightarrow R$ }
{4}	(4) $A \rightarrow \neg H$	{Rule P}
{4}	(5) $H \rightarrow \neg A$	{Rule T i, e $P \rightarrow Q \Leftrightarrow \neg Q \rightarrow \neg P$ }
{3,5}	(6) $E \rightarrow A$	{Rule T i, e $P \rightarrow Q, Q \rightarrow R \Rightarrow P \rightarrow R$ }
{6}	(7) $\neg E \vee \neg A$	{Rule T i, e $P \rightarrow Q \Rightarrow \neg P \vee Q$ }
{7}	(8) $\neg(E \wedge A)$	{By Demorgans properties i, e $\neg E \vee \neg A \Rightarrow \neg(E \wedge A)$ }
{9}	(9) $E \wedge A$	{Rule P}
{8,9}	(10) $\neg(E \wedge A) \wedge (E \wedge A)$	{Rule T i, e $P, Q \Rightarrow P \wedge Q$ }

There is conjunction implies a contradiction (FALSE); hence the given premises are inconsistent.

OR

(C) Define conjunctive normal form and obtain a conjunctive normal form of

$$\neg(P \vee Q) \Leftrightarrow (P \wedge Q)$$

**Conjunctive Normal Form:-** Any formula which is equivalent to a given formula and which consist of product of elementary sum is called conjunctive normal form of given formula.

**Proof:-**

$$\begin{aligned} & \neg(P \vee Q) \Leftrightarrow (P \wedge Q) \\ & \text{by } R \Leftrightarrow S \Leftrightarrow (R \rightarrow S) \wedge (S \rightarrow R) \\ & \Leftrightarrow [\neg(P \vee Q) \rightarrow (P \wedge Q)] \wedge [(P \wedge Q) \rightarrow \neg(P \vee Q)] \\ & \Leftrightarrow [\neg\neg(P \vee Q) \vee (P \wedge Q)] \wedge [\neg(P \wedge Q) \vee \neg(P \vee Q)] \quad \{P \rightarrow Q \Rightarrow \neg P \vee Q\} \\ & \Leftrightarrow [(P \vee Q) \vee (P \wedge Q)] \wedge [\neg(P \vee \neg Q) \vee \neg(P \vee Q)] \quad \{By Demorgans property \& \neg\neg P \Rightarrow P\} \\ & \Leftrightarrow [(P \vee Q \vee P) \wedge (P \vee Q \vee Q)] \wedge [(\neg P \vee \neg Q \vee \neg P) \wedge (\neg P \vee \neg Q \vee \neg Q)] \quad \{By Distributive property\} \\ & \Leftrightarrow ((P \vee P) \vee Q) \wedge ((Q \vee Q) \vee P) \wedge ((\neg P \vee \neg P) \vee \neg Q) \wedge ((\neg Q \vee \neg Q) \vee \neg P) \quad \{By Associative property\} \\ & \Leftrightarrow (P \vee Q) \wedge (Q \vee P) \wedge (\neg P \vee \neg P) \wedge (\neg Q \vee \neg P) \quad \{P \vee P = P\} \\ & \Leftrightarrow (P \vee Q) \wedge (P \vee Q) \wedge (\neg P \vee \neg P) \wedge (\neg P \vee \neg Q) \quad \{By commulative property\} \\ & \Leftrightarrow (P \vee Q) \wedge (\neg P \vee \neg P) \quad \{P \wedge P = P\} \end{aligned}$$

It is the form of product of elementary sum of min terms.

Hence it is form of Principal Conjunction Normal Form.

(D) Show that  $R \wedge (P \vee Q)$  is a valid conclusion from the premises  $P \vee Q$ ,

$Q \rightarrow R, P \rightarrow M$  and  $\neg M$ .

**Solution:**

{1}	(1) $\neg M$	{Rule P}
{2}	(2) $P \rightarrow M$	{Rule P}
{2}	(3) $\neg M \rightarrow \neg P$	{Rule T : $P \rightarrow Q \Leftrightarrow \neg Q \rightarrow \neg P$ }
{3}	(4) $\neg P$	{Rule T : $P, P \rightarrow Q \Leftrightarrow Q$ }

{5}	(5) $P \vee Q$	{ Rule P
{5}	(6) $\neg P \rightarrow Q$	{Rule T : $\neg P \rightarrow Q \Leftrightarrow \neg \neg P \vee Q \Leftrightarrow P \vee Q$
{4,6}	(7) Q	{ Rule T : $P, P \rightarrow Q \Leftrightarrow Q$
{8}	(8) $Q \rightarrow R$	{ Rule P
{7,8}	(9) R	{ Rule T : $Q, Q \rightarrow R \Leftrightarrow R$
{9,5}	(10) $R \wedge (P \vee Q)$	{Rule T : $P, Q \Leftrightarrow P \wedge Q$

Hence Proved.

### 3. EITHER

(A) Find an explicit formula for the sequence defined by

$$a_n = 4a_{n-1} + 5a_{n-2} \text{ with initial conditions } a_1 = 2, a_2 = 6.$$

**Solution:**

First find sequence for recurrence relation

$$a_n = 4a_{n-1} + 5a_{n-2}$$

$$\text{For } n = 3 \quad a_3 = 4a_{3-1} + 5a_{3-2}$$

$$= 4a_2 + 5a_1$$

$$= 4(6) + 5(2)$$

$$= 24 + 10$$

$$= 34$$

$$\text{For } n = 4 \quad a_4 = 4a_{4-1} + 5a_{4-2}$$

$$= 4a_3 + 5a_2$$

$$= 4(34) + 5(6)$$

$$= 166$$

$$\text{For } n = 5 \quad a_5 = 4a_{5-1} + 5a_{5-2}$$

$$= 4a_4 + 5a_3$$

$$= 4(166) + 5(34)$$

$$= 834$$

$\therefore$  Sequence is 2,6,34,166,834----

The recurrence relation  $a_n = 4a_{n-1} + 5a_{n-2}$  is linear homogeneous

Equation of degree 2.

It associated equation is

$$x^2 = 4x + 5$$

Rewriting this as

$$x^2 - 4x - 5 = 0$$

$$x^2 - 5x + x - 5 = 0$$

$$(x-5)(x+1) = 0$$

$$x = 5 \text{ or } x = -1$$

The roots of the equation is  $s_1 = 5$  and  $s_2 = -1$

Now, by theorem(i)

We can find value of u and v

$$\text{From } a_n = us_1^n + vs_2^n \quad \text{-----(A)}$$

For  $n=1$

$$a_1 = us_1 + vs_2$$

$$2 = u(5) + v(-1)$$

$$2 = 5u - v \quad \text{-----(i)}$$

For  $n = 2$

$$a_2 = us_1^2 + vs_2^2$$

$$6 = u(5)^2 + v(-1)^2$$

$$6 = 25u + v \quad \text{-----(ii)}$$

Solving equation (i) and (ii)

$$5u - v = 2$$

$$25u + v = 6$$

$$+ \quad + \quad +$$

---

$$30u = 8$$

$$\boxed{u = 8/30}$$

Putting values of u in equation (i)

$$2 = 5(8/30) - v$$

$$2 = (8/6) - v$$

$$2 = 8/6 - 2$$

$$2 = 8 - 12/6$$

$$v = -4/6$$

$$\boxed{v = -2/3}$$

Put value of  $u_1, v_1, s_1$  and  $s_2$  in equation (A)

$$a_n = us_1^n + vs_2^n$$



$$a_n = (8/30) (5)^n + (-2/3) (-1)^n$$

$$a_n = 8/30 (5)^n + -2/3 (-1)^n$$

∴ Which is required formula.

**(B)** Let  $A = \{a, b, c, d\}$  and let  $R$  be the relation on  $A$  that has the matrix.

$$M_r = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 \end{bmatrix}$$

Draw the diagram of  $R$ . Find the indegree and outdegree of all vertices.

OR

**(C)** Let  $A = \mathbb{Z}$  and let  $R = \{(a, b) \in A \times A : a \equiv r \pmod{2} \text{ and}$

$$b \equiv r \pmod{2}\}.$$

Show that the relation  $R$  is an equivalence relation.

**(D)** Prove that,

Let  $R$  be a relation on a set  $A$ . Then  $R^\infty$  is the transitive

Closure of  $R$ .

**Solution:-**

If  $a$  &  $b$  are in the set  $A$  then  $a R^\infty b$  iff there is a path in  $R$  from  $a$  to  $b$ .

Now  $R^\infty$  is certainly transitive if  $a R^\infty b$  and  $b R^\infty c$  then the composition of path from  $a$  to  $b$  and  $b$  to  $c$  from sub path from  $a$  to  $b$  in  $R$  and so  $a R^\infty c$ . To show that  $R^\infty$  is the smallest transitive relation containing  $R$ .

We must show that if  $S$  is any transitive relation on  $A$  and  $R \subseteq S$  then  $R^\infty$  is smallest of  $S$  ( $R^\infty \subseteq S$ )

We know that if  $S$  is transitive then,

$S^n \circ S^n \subseteq S$  for all  $n$ . i.e if  $a$  &  $b$  are connected by a path of length  $n$  then  $a S b$  it follows that

$$S^\infty = \bigcup_{n \neq 1}^{\infty} a S^n b$$

$$S^n \text{ OR } S^\infty = \bigcup_{n \neq 1}^{\infty} S^n \subseteq S$$

It is also true if  $R \subseteq S$  then  $R^\infty$  is subset of  $S^\infty$  since any path in  $R$  is also path in  $S$  putting this fact together, we see that,

If  $R \subseteq S$  and  $S$  is transitive on  $A$  then,

$$R^\infty \text{ subset of } S^\infty \subseteq S, \quad R^\infty \subseteq S$$

This means that  $R$  is the smallest of all transitive relation on  $A$ , that contains  $R^\infty$ .

Hence theorem is proved.

#### 4. EITHER

(A) Show that if  $n$  is a positive integer and  $P^2 | n$ , where  $p$  is prime number, then  $D_n$  is not Boolean algebra.

**Proof:-**

Suppose that  $P^2 | n$ ,  $a/b \Rightarrow b = ac$  for sum  $c$

So,  $n = p^2 \cdot q$  for some positive integer  $q$ .

$\because p$  is also a divisor of  $n$  and  $p$  is an element of  $D_n$ , (Divisor of  $n$  number i.e  $D_n$ )

$\rightarrow$  natural number i.e if  $D_n$  is a Boolean Algebra. Then  $p$  must have a complement  $p'$ .

Then  $GCD(p, p') = 1$

And  $LCM(p, p') = 1$

$GCD.LCM(p, p') = p.p'$

i.e  $p.p' = n$

so,  $p' = n/p = 1$

$p' = p.q$

i.e  $GCD(p, p.q) = 1$

This is impossible.

$\because p$  &  $p.q$  have  $p$  as a common divisor.

(B) Find the Hamiltonian circuit for the given graph.

**Answer:**

**Hamiltonian Graph:**

A Hamiltonian graph is a graph that on a Hamiltonian path.

A Hamiltonian path uses each vertex exactly once but edges be include.

**OR**

(C) Let  $L$  is a bounded distributive lattice. Prove that, if a complement of an element in  $L$  exists then it is unique.

**Proof:-**

Lattice suppose  $a'$  and  $a''$  are two complements of an element  $a \in L$

$$a \wedge a' = 0 \quad a \wedge a'' = 0$$

$$a \vee a' = I \quad a \vee a'' = I$$

We show that  $a' = a''$

Now

$$\begin{aligned}
a' &= a' \wedge I \\
a' &= a' \wedge (a \vee a'') \\
a' &= (a' \wedge a) \vee (a' \wedge a'') \\
a' &= 0 \vee (a' \wedge a'') \\
a' &= a' \wedge a'' \quad \text{----- (1)}
\end{aligned}$$

And

$$\begin{aligned}
a'' &= a'' \wedge I \\
a'' &= a'' \wedge (a \vee a') \\
a'' &= (a'' \wedge a) \vee (a'' \wedge a') \\
a'' &= 0 \vee (a'' \wedge a') \\
a'' &= a'' \wedge a' \\
a'' &= a' \wedge a'' \quad \{by\ commutative\ property} \quad \text{----- (2)}
\end{aligned}$$

$\therefore a' = a''$  (From equation (1) and (2) )

$\therefore$  Complement if exists is unique.

$\therefore$  *proved.*

**(D) Prove that :-**

**A tree with n vertices has n-1 edges.**

**Proof:-**

Consider tree  $T(V,E)$

By using mathematical induction on the number of vertices, n in T.

Suppose, it is true  $n=m(\geq 2)$ .

m- some positive integer.

To prove for  $n = m+1$

Suppose, T has  $m+1$  vertices.

If we remove an edge with end points u&v from T.

Then we are left with two sub trees  $T_1(V_1,E_1)$  and  $T_2(V_2,E_2)$

Such that  $|V| = |V_1| + |V_2|$  AND  $|E| = |E_1| + |E_2|$

$T_1$  &  $T_2$  are connected with number cycles and having vertices less than n,

i.e  $|V_1| \leq n$  &  $|V_2| \leq n$ ,

i.e  $|E_1| = |V_1| - 1$ ;  $|E_2| = |V_2| - 1$

$|V| = |V_1| + |V_2|$

$|V| = (|E_1| + 1) + (|E_2| + 1)$

$|V| = (|E_1| + |E_2| + 1) + 1$

$|V| = |E_1| + 1$

OR

$|E| = (|V_1| - 1) + (|V_2| - 1) + 1$

$|E| = (|V_1| + |V_2| - 1)$

$$|E| = m+1-1$$

$$|E| = m$$

This is prove that T has m edges which is sequence number.

## 5. EITHER

(A) Define :-

- (i) **Semigroup**
- (ii) **Monoid**
- (iii) **Subsemigroup**
- (iv) **Group homomorphism.**

Answer:-

(i) **Semigroup:-**

Let S be a non-empty set and \* be a binary operation on S. The algebraic system (S, \*) is called a semigroup if the operation \* is

(1) The operation \* is a closed operation on set A.

(2) The operation \* is an associative operation.

Or (S, \*) is a semigroup if for any  $x, y, z \in S$ ,

$$(x * y) * z = x * (y * z)$$

**Free semigroup:**

If \* is an associative binary operation, and (A, \*) is a semigroup. The semigroup(A, \*) is called free semigroup by A.

Ex:

Consider an algebraic system (S, \*) where  $S = \{1,2,3,5,7,9, \dots\}$  the set of all positive odd integers and \* is a binary operation means multiplication. Determine whether (S, \*) is a semigroup.

(ii) **Monoid:-**

Let us consider an algebraic system (M, \*), where \* is a binary operation on M. Then the system (M, \*) is said to be a monoid if it satisfies the following properties:

- (1) The operation \* is a closure operation on set A.
- (2) The operation \* is an associative operation.
- (3) There exists an identify element w. r. t. The operation \*.

Ex:-

Consider an algebra system (N, +), where the set  $N = \{0,1,2,3, \dots\}$  the set of natural numbers and + is an addition operation. Determine whether (N,+) is a monoid.

(iii) **Subsemigroup:-**

Let (S, \*) be a semigroup and  $T \subseteq S$ , if the set T is closed under the operation \* then (T, \*) is said to be subsemigroups of (S, \*).

Ex:

Consider a semigroup (N,+), where N is the set of all natural number and + is an addition operation.

The algebraic system  $(E,+)$  is a subsemigroup of  $(N,+)$ , where E is a set of all +ve even integer.

**(iv) Group homomorphism:-**

Let  $(S,*)$  and  $(T,*)$  be two semigroups. An everywhere defined function  $f: S \rightarrow T$  is called a homomorphism from  $(S,*)$  and  $(T,*)$

$$\text{If } (a * b) = f(a) * f(b)$$

For all a and b in S.

If f is also onto.

We say that T is a homomorphic image of S.

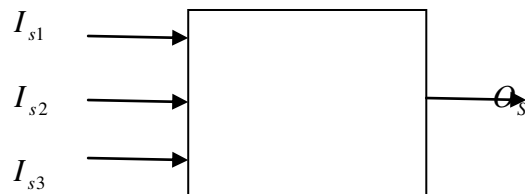
**(B) Define finite state machine. Construct digraph of machine whose table is**

	a	b	c
S0	S0	S0	S0
S1	S2	S3	S2
S2	S1	S0	S3
S3	S3	S2	S3

Answer:-

**Finite state machines:**

Finite state machine that accepts more than one input and gives single output then it is Finite-State-Machine.



**Fig:** Finite-State Machine.

Where  $I_s$  = input signal,  $O_s$  = output signal

Here  $I_{s1}, I_{s2}, I_{s3}, \dots, I_{sn}$  is number of input which gives single output signal as shown in fig.

**Definition:**

Finite-State Machine define by 3-tuple (triple)

$$M = (S, I, F)$$

Where S = finite set of state of machine i,e  $S = \{S_0, S_1, \dots, S_n\}$

I = finite set of input of machines.

F = Is the state transition function i,e  $F = \{f_x / x \in I\}$

(for each  $x \in I$ , a function  $f_x : S \rightarrow S$ )

**Solution:**

Let finite-state machines define by 3-tuple.

$$M = (S, I, F)$$

Where  $S = \{S_0, S_1, S_2, S_3\}$

$$I = \{a, b, c\}$$

F = state transition function.

**Given that:** Transition Table

	a	b	c
S0	S0	S0	S0
S1	S2	S3	S2
S2	S1	S0	S3
S3	S3	S2	S3

From the above transition table draw digraph for the machine as follows.

**Digraph:**

**OR**

**(C) Let G be the set of all non zero real numbers and Let  $a * b = ab/2$  show that  $(G, *)$  is an abelian group.**

**Solution:-**

If a, b are element in G the  $\frac{ab}{2}$  is a non-zero real number.

**To show:**  $(G, *)$  is an abelian group.

**Closure property:**

The set G is closed under the operation \*.

Since,  $a*b = \frac{ab}{2}$  is a real number.

Hence, belongs to G.

**Associative property:**

The operation \* is associative.

Let a, b, c  $\in$  G, then

We have

$$\begin{aligned} (a * b) * c &= \left( \frac{ab}{2} \right) * c \\ &= \frac{(ab)c}{4} \end{aligned}$$

$$= \frac{abc}{4}$$

Similarly,  $a * (b * c) * a = \left( \frac{ab}{2} \right)$

$$= \frac{a(bc)}{4}$$

$$= \frac{abc}{4}$$

**Identity :**

To find the identity element.

Suppose that 'e' is a +ve real number.

Then,  $e * a = a$ , where  $a \in G$

$$\frac{ea}{2} = a \quad \text{or} \quad e = 2$$

Similarly,  $a * e = a$

$$\frac{ae}{2} = a \quad \text{or} \quad e = 2$$

Thus, the identity element in G is 2.

**Inverse :**

Suppose that  $a \in G$ .

If  $a^{-1} \in G$  is an inverse of a, then  $a * a^{-1} = 2$ .

Therefore,  $\frac{aa^{-1}}{2} = 2 \quad \text{or} \quad a^{-1} = \frac{4}{a}$

Thus, the inverse of element 'a' in G is  $\frac{4}{a}$

**Commutative :**

The operation \* on G is commutative.

Since,  $a * b = \frac{ab}{2} = b * a$

Thus, the algebraic system  $(G, *)$  is closed, associative, identity element, inverse and commutative.

Hence, the system  $(G, *)$  is an abelian group.

**(D) consider the semigroup  $(\mathbb{Z}, +)$  and the equivalence relation R on Z defined by  $aRb$  if and only if  $a \equiv b \pmod{2}$ . Show that this relation is a congruence relation.**

**Solution:**

Remember that if  $a \equiv b \pmod{2}$ , then  $2 \mid a - b$ .

We now show that this relation is a congruence relation as follows.

$$a \equiv b \pmod{2}$$

and  $c \equiv d \pmod{2}$

then  $2$  divide  $a - b$

and  $2$  divide  $c - d$

so  $a - b = 2m$

and  $c - d = 2n$

where  $m$  and  $n$  are  $\mathbb{Z}$ .

Adding, we have

$$(a - b) + (c - d) = 2m + 2n$$

Or  $(a + c) - (b + d) = 2(m + n)$

So,  $a + c \equiv b + d \pmod{2}$

Hence, the relation is congruence relation.



**EITHER**

**(A) Define:-**

- i. Boolean Matrix**
- ii. Join of Boolean Matrices**
- iii. Meet of Boolean Matrices**
- iv. Boolean product and**

Show that  $A \ominus (B \ominus C) = (A \ominus B) \ominus C$

$$\text{If } A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \end{bmatrix} \quad B = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 0 & 1 \\ 1 & 0 & 1 \end{bmatrix} \quad C = \begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 1 \\ 0 & 0 & 1 \end{bmatrix}$$

**Solution:**

Let us assume that

$$A = [a_{ij}] m \times n$$

$$B = [b_{jk}] m \times n$$

$$C = [c_{kl}] m \times n$$

$$(A \vee B) \vee C = A \vee (B \vee C)$$

i.e.  $(A+B)+C = A+(B+C)$

Now,

$$A + B = [a_{ij}] m \times n + [b_{jk}] m \times n$$

$$= [a_{ij} + b_{jk}] m \times n$$

$$(A + B) + C = [a_{ij} + b_{jk}] m \times n + [c_{kl}] m \times n$$

$$= [a_{ij} + b_{jk} + c_{kl}] m \times n \longrightarrow (1)$$

$$(B + C) = [b_{jk}] m \times n + [c_{kl}] m \times n$$

$$= [b_{jk} + c_{kl}] m \times n$$

$$(B + C) + A = [a_{ij}] m \times n + [b_{jk} + c_{kl}] m \times n$$

$$= [a_{ij} + b_{jk} + c_{kl}] m \times n \longrightarrow (2)$$

From equation (1) & (2)

We get  $(A \vee B) \vee C = A \vee (B \vee C)$

i.e.

$$A \ominus (B \ominus C) = (A \ominus B) \ominus C$$

**(B) Make truth table for**

**i.  $(p \wedge q) \vee (\neg p)$**

ii.  $(p \downarrow q) \downarrow r$

Solution:  $(p \wedge q) \vee (\neg p)$

Truth Table:

p	q	$(p \wedge q)$	$(\neg p)$	$(p \wedge q) \vee (\neg p)$
T	T	T	F	T
T	F	F	F	F
F	T	F	T	T
F	F	F	T	T

$(p \downarrow q) \downarrow r$

Truth Table:

p	q	r	$(p \downarrow q)$	$(p \downarrow q) \downarrow r$
T	T	T	F	F
T	T	F	F	T
T	F	T	F	F
T	F	F	F	T
F	T	T	F	F
F	T	F	F	T
F	F	T	T	F
F	F	F	T	F

OR

(C) Use Induction method to prove that

$$\overline{(\bigcap_{i=1}^n A_i)} = \bigcup_{i=1}^n \overline{A_i}$$

Solution:

$$\overline{A_1 \cap A_2 \cap A_3 \dots \cap A_n} = \overline{A_1} \cup \overline{A_2} \cup \overline{A_3} \dots \cup \overline{A_n}$$

Basic Steps:

Let  $n=1$

$$P(A) = \overline{A_1} = \overline{A_1}$$

$P(n)$  is true for  $n=1$

Induction Steps:

Let us assume that  $P(n)$  is true for  $n=k$

$$P(k) = \overline{(\bigcap_{i=1}^k A_i)} = \bigcup_{i=1}^k \overline{A_i}$$

$$\therefore \overline{A_1 \cap A_2 \cap A_3 \dots \cap A_k} = \overline{A_1} \cup \overline{A_2} \cup \overline{A_3} \dots \cup \overline{A_k} \longrightarrow (1)$$

Now,

We have to prove that P(n) is true for n=k+1

$$P(k+1) = \overline{\left(\bigcap_{i=1}^{k+1} A_i\right)} = \bigcup_{i=1}^{k+1} \overline{A_i}$$

$$\overline{A_1 \cap A_2 \cap A_3 \dots \cap A_{k+1}} = \overline{A_1} \cup \overline{A_2} \cup \overline{A_3} \dots \cup \overline{A_{k+1}}$$

L.H.S.

$$\Rightarrow \overline{A_1 \cap A_2 \cap A_3 \dots \cap A_{k+1}}$$

$$\Rightarrow \overline{A_1 \cap A_2 \cap A_3 \dots A_k \cap A_{k+1}}$$

$$\Rightarrow \overline{A_1 \cap A_2 \cap A_3 \dots A_k} \cup \overline{A_{k+1}} \quad \{\text{By demorgans property } \overline{A \cap B} = \overline{A} \cup \overline{B}\}$$

$$\Rightarrow \overline{A_1} \cup \overline{A_2} \cup \overline{A_3} \dots \cup \overline{A_{k+1}} \quad \{\text{From eq. (1)}\}$$

$\Rightarrow$  R.H.S.

Hence Proved

P(n) is true for n=k+1

**(D)** Let m and n be integers. Prove that  $n^2 = m^2$  if and only if  $m=n$  or  $m=-n$ . Also prove that  $3|(n^3 - n)$  for every positive integers

**EITHER**

**Q.2**

**(A) Show that**

$$(P \vee Q) \wedge \neg(P \wedge (\neg Q \vee \neg R)) \vee (\neg P \wedge \neg Q) \vee \neg(P \wedge \neg R)$$

**Is tautology without using truth table**

$$\Rightarrow [(P \vee Q) \wedge \neg(P \wedge (\neg Q \vee \neg R))] \vee \neg(P \wedge \neg Q) \vee \neg(P \wedge \neg R) \quad \{\text{by associative property}\}$$

$$\Rightarrow [(P \vee Q) \wedge \neg(P \wedge \neg(Q \vee R))] \vee \neg(P \wedge \neg Q) \vee \neg(P \wedge \neg R) \quad \{\text{By demorgans and distributive prop. Respt.}\}$$

$$\Rightarrow [(P \vee Q) \wedge \neg(P \vee (Q \wedge R))] \vee \neg(P \wedge \neg(Q \wedge R)) \quad \{\text{by demorgans property}\}$$

$$\Rightarrow [(P \vee Q) \wedge (P \vee (Q \wedge R))] \vee \neg(P \vee (Q \wedge R)) \quad \{\text{by demorgans property } \& \neg \neg A \Rightarrow A\}$$

A

$$\Rightarrow [P \vee (Q \wedge (Q \wedge R))] \vee \neg(P \vee (Q \wedge R)) \quad \{\text{by distributive property.}\}$$

$$\Rightarrow [P \vee (Q \wedge Q) \wedge R] \vee \neg(P \vee (Q \wedge R)) \quad \{\text{by associative property}\}$$

$$\Rightarrow (P \vee (Q \wedge R)) \vee \neg(P \vee (Q \wedge R)) \quad \{\text{by idempotent property } \& P \wedge P \Rightarrow P\}$$

$$\Rightarrow T \quad \{P \vee \neg P \Rightarrow T\}$$

Hence Proved

**(B) Define conjunctive normal form and obtain a conjunctive normal form of**

$$\neg(P \vee Q) \leftrightarrow (P \wedge Q)$$

**Conjunctive Normal Form:-** Any formula which is equivalent to a given formula and which consist of product of elementary sum is called conjunctive normal form of given formula.

**Proof:-**

$$\neg(P \vee Q) \leftrightarrow (P \wedge Q)$$

$$\text{by } R \leftrightarrow S \Leftrightarrow (R \rightarrow S) \wedge (S \rightarrow R)$$

$$\Leftrightarrow [\neg(P \vee Q) \rightarrow (P \wedge Q)] \wedge [(P \wedge Q) \rightarrow \neg(P \vee Q)]$$

$$\Leftrightarrow [\neg\neg(P \vee Q) \vee (P \wedge Q)] \wedge [\neg(P \wedge Q) \vee \neg(P \vee Q)] \quad \{P \rightarrow Q \Rightarrow \neg P \vee Q\}$$

$$\Leftrightarrow [(P \vee Q) \vee (P \wedge Q)] \wedge [(\neg P \vee \neg Q) \vee \neg(P \vee Q)] \quad \{\text{By Demorgans property \& } \neg\neg P \Rightarrow P\}$$

$$\Leftrightarrow [(P \vee Q \vee P) \wedge (P \vee Q \vee Q)] \wedge [(\neg P \vee \neg Q \vee \neg P) \wedge (\neg P \vee \neg Q \vee \neg Q)] \quad \{\text{By Distributive property}\}$$

$$\Leftrightarrow ((P \vee P) \vee Q) \wedge ((Q \vee Q) \vee P) \wedge ((\neg P \vee \neg P) \vee \neg Q) \wedge ((\neg Q \vee \neg Q) \vee \neg P) \quad \{\text{By Associative property}\}$$

$$\Leftrightarrow (P \vee Q) \wedge (Q \vee P) \wedge (\neg P \vee \neg Q) \wedge (\neg Q \vee \neg P) \quad \{P \vee P = P\}$$

$$\Leftrightarrow (P \vee Q) \wedge (P \vee Q) \wedge (\neg P \vee \neg Q) \wedge (\neg P \vee \neg Q) \quad \{\text{By commulative property}\}$$

$$\Leftrightarrow (P \vee Q) \wedge (\neg P \vee \neg Q) \quad \{P \wedge P = P\}$$

It is the form of product of elementary sum of min terms.

Hence it is form of Principal Conjunction Normal Form.

**(C) Show that  $R \wedge (P \vee Q)$  is a valid conclusion from the premises  $P \vee Q$ ,  $Q \rightarrow R$ ,  $P \rightarrow M$  and  $\neg M$ .**

<b>Solution:</b>	{1}	(1) $\neg M$	{Rule P
	{2}	(2) $P \rightarrow M$	{ Rule P
$\neg A \quad \pi$		(3) $\neg M - \neg P$	{ Rule T : $P \rightarrow Q \Leftrightarrow \neg Q - \neg P$
	{3}	(4) $\neg P$	{ Rule T : $P \rightarrow Q \Leftrightarrow Q$
		(5) $P \vee Q$	{ Rule P
$= \quad \neg^2$		(6) $\neg P \rightarrow Q$	{Rule T : $\neg P \rightarrow Q \Leftrightarrow \neg\neg P \vee Q \Leftrightarrow P \vee Q$
	{4,6}	(7) Q	{ Rule T : $P \rightarrow Q \Leftrightarrow Q$
	{8}	(8) $Q \rightarrow R$	{ Rule P
	{7,8}	(9) R	{ $Q, Q \rightarrow R \Leftrightarrow R$
	{9,5}	(10) $R \wedge (P \vee Q)$	{Rule T : $P, Q \Leftrightarrow P \wedge Q$

Hence Proved.

**(D) Show that**

$$(x)(P(x) \rightarrow Q(x)) \wedge (x)(Q(x) \rightarrow R(x)) \Rightarrow (x)(P(x) \rightarrow R(x))$$

**Solution:**

Given Premises are

$$(x) (P(x) \rightarrow Q(x)) \wedge (x) (Q(x) \rightarrow R(x))$$

We have to derive,

$$(x) (P(x) \rightarrow R(x))$$

{1}	(1) $(x) (P(x) \rightarrow Q(x))$	{Rule P}
{2}	(2) $P(y) \rightarrow Q(y)$	{Rule US : $(x)A(x) \rightarrow A(y)$ }
{3}	(3) $(x) (Q(x) \rightarrow R(x))$	{Rule P}
{3}	(4) $Q(y) \rightarrow R(y)$	{Rule US : $(x)A(x) \rightarrow A(y)$ }
{2,4}	(5) $P(y) \rightarrow R(y)$	{Rule T : $P \rightarrow Q, Q \rightarrow R \Rightarrow P \rightarrow R$ }
{5}	(6) $(x) (P(x) \rightarrow R(x))$	{Rule UG : $A(y) \rightarrow_x A(x)$ }

Hence Prove.

**Q.3 EITHER**

(A) Define Cartesian product of two sets, partition of a set and prove that  $A \times (B \cup C) = (A \times B) \cup (A \times C)$

**Solution:**

Cartesian product of two sets:

If A and B are the two non-empty sets, we define the product set or Cartesian product  $A \times B$  as the set of all ordered pair  $(a,b)$  with  $a \in A$  and  $b \in B$ .

Thus

$$A \times B = \{(a,b) | a \in A \text{ and } b \in B\}$$

Ex: let  $A = \{1,2,3\}$  and  $B = \{r,s\}$ . Determine the product set of  $A \times B$  and  $B \times A$ .

Solution: Let  $A = \{1,2,3\}$  and  $B = \{r,s\}$

To find : (1)  $A \times B$  (2)  $B \times A$

(1) The Cartesian product of A and B is  
 $A \times B = \{(1,r), (1,s), (2,r), (2,s), (3,r), (3,s)\}$

(2) The Cartesian product or product sets of B and A is  
 $B \times A = \{(r,1), (r,2), (r,3), (s,1), (s,2), (s,3)\}$

**Prove that  $A \times (B \cup C) = (A \times B) \cup (A \times C)$**

Solution : Let  $(x,y) \in A \times (B \cup C) \Rightarrow x \in A$  and  $y \in B \cup C$

$$\Rightarrow x \in A \text{ and } (y \in B \text{ or } y \in C)$$

$$\Rightarrow (x \in A \text{ and } y \in B) \text{ or } (x \in A \text{ and } y \in C)$$

$$\Rightarrow (x,y) \in A \times B \text{ or } (x,y) \in A \times C$$

$$\Rightarrow (x,y) \in (A \times B) \cup (A \times C)$$

$$\text{Therefore, } A \times (B \cup C) \subset (A \times B) \cup (A \times C) \dots \dots (1)$$

Now,

Conversely

$$\text{Let } (x,y) \in (A \times B) \cup (A \times C)$$

$$\Rightarrow (x,y) \in (A \times B) \text{ or } (x,y) \in (A \times C)$$

$$\Rightarrow (x \in A \text{ and } y \in B) \text{ or } (x \in A \text{ and } y \in C)$$

$$\Rightarrow x \in A \text{ and } y \in B \text{ or } B \text{ or } y \in C$$

$$\Rightarrow x \in A \text{ and } y \in (B \cup C)$$

$$\text{Therefore, } (A \times B) \cup (A \times C) \subset A \times (B \cup C) \dots \dots (2)$$

From (1) and (2), we have

$$A \times (B \cup C) = (A \times B) \cup (A \times C)$$

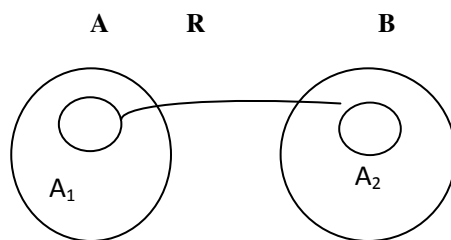
Hence proved.

**(B) Let R is a relation from A to B, and let A<sub>1</sub> and A<sub>2</sub> be subsets of A. Then show that**

$$(1) R(A_1 \cup A_2) = R(A_1) \cup R(A_2) \text{ and}$$

$$(2) R(A_1 \cap A_2) = R(A_1) \cap R(A_2)$$

**Solution:**



$$(1) R(A_1 \cup A_2) = R(A_1) \cup R(A_2)$$

$$\text{Let } y \in R(A_1 \cup A_2)$$

$$\Rightarrow \exists x \in A_1 \cup A_2 \text{ s.t. } (x,y) \in R \text{ or } xRy$$

$$\Rightarrow \exists x \in A_1 \text{ or } x \in A_2 \text{ or } x \in A_1 \& A_2 \text{ s.t. } (x,y) \in R$$

$$\text{If } x \in A_1 \text{ s.t. } (x,y) \in R \text{ then } y \in R(A_1)$$

$$\text{And If } x \in A_2 \text{ s.t. } (x,y) \in R \text{ then } y \in R(A_2)$$

Now,

$$y \in R(A_1) \text{ or } y \in R(A_2) \Rightarrow y \in R(A_1) \cup R(A_2)$$

$$\therefore R(A_1 \cup A_2) \subseteq R(A_1) \cup R(A_2) \dots\dots(1)$$

Now, suppose,

$$y \in R(A_1) \cup R(A_2)$$

$$\Rightarrow y \in R(A_1) \text{ or } y \in R(A_2) \text{ or } y \in R(A_1) \cap R(A_2)$$

$$\text{If } y \in R(A_1) \text{ then } \exists x \in A_1 \text{ s.t. } (x,y) \in R$$

$$\text{If } y \in R(A_2) \text{ then } \exists x \in A_2 \text{ s.t. } (x,y) \in R$$

$\therefore$  we have

$$\exists x \in A_1 \text{ or } A_2 \text{ s.t. } (x,y) \in R$$

$$\text{i.e. } \exists x \in A_1 \cup A_2 \text{ s.t. } (x,y) \in R \Rightarrow y \in R(A_1 \cup A_2)$$

$$\therefore R(A_1) \cup R(A_2) \subseteq R(A_1 \cup A_2) \dots\dots(2)$$

From (1) & (2) we get

$$R(A_1 \cup A_2) = R(A_1) \cup R(A_2)$$

Hence proved

$$(2) R(A_1 \cap A_2) = R(A_1) \cap R(A_2)$$

$$\text{Let } y \in R(A_1 \cap A_2)$$

$$\Rightarrow \exists x \in A_1 \cap A_2 \text{ s.t. } (x,y) \in R$$

$$\Rightarrow \exists x \in A_1 \text{ and } x \in A_2 \text{ s.t. } (x,y) \in R$$

Now,

$$x \in A_1 \text{ and } (x,y) \in R \Rightarrow y \in R(A_1) \text{ and } x \in A_2 \text{ and } (x,y) \in R \Rightarrow y \in R(A_2)$$

$$\therefore y \in R(A_1) \cap R(A_2) \Rightarrow y \in R(A_1) \cap R(A_2)$$

$$R(A_1 \cap A_2) = R(A_1) \cap R(A_2)$$

(C) Let  $a = \{1,2,3\}$  and let the relation **R** and **S** on **A** are

$$R = \{(1,1), (1,2), (2,1), (1,3), (3,1)\}$$

$$S = \{(1,1), (1,2), (2,1), (2,2), (3,3)\}$$

Find  $\bar{R}, R^{-1}, \bar{S}, S^{-1}, S \cap S, R \cup S$

(D) Let A be set with  $|A| = n$  and Let R be a relation on A then prove that  $R^\infty = R \cup R^2 \cup \dots \cup R^n$

#### Q.4 EITHER

(A) Let the number of edges of G be M. Then prove that G has a Hamiltonian circuit if

$$m \geq \frac{1}{2} (n^2 - 3n + 6)$$

**Proof:**

Suppose U & V are two vertices of graph G that are not adjacent.

Let H be a graph product by elementary vertices U & K from G.

$\therefore$  H has  $n-2$  vertices

$\therefore$  no. of edges in H are  $m - \text{deg}(U) - \text{deg}(V)$ .

$\therefore$  maximum no. of edges in H are  $(n-2)!$

$$\frac{(n-2)!}{2!(n-2-2)!}$$

$$\Rightarrow \frac{(n-2)!}{2!(n-4)!}$$

$$\Rightarrow \frac{(n-2)(n-3)(n-4)!}{2!(n-4)!}$$

$(n-4)!$  Get Cancel

$$\Rightarrow \frac{(n-2)(n-3)}{2}$$

$$\Rightarrow \frac{1}{2} (n^2 - 5n + 6)$$

$$\therefore m - \text{deg}(U) - \text{deg}(V) \leq \frac{1}{2} (n^2 - 5n + 6)$$

$$m - \frac{1}{2} (n^2 - 5n + 6) \leq m - \text{deg}(U) - \text{deg}(V)$$

or

$$\text{deg}(U) + \text{deg}(V) \geq m - \frac{1}{2} (n^2 - 5n + 6)$$

$$\text{deg}(U) + \text{deg}(V) \geq \frac{1}{2} (n^2 - 3n + 6) - \frac{1}{2} (n^2 - 5n + 6)$$

$$\text{deg}(U) + \text{deg}(V) \geq \frac{1}{2} [n^2 - 3n + 6 - n^2 + 5n - 6]$$



$n^2$  &  $-n^2$  and  $+6$  &  $-6$  get cancel

$$\deg(U) + \deg(V) \geq \frac{1}{2} \times 2n$$

$$\deg(U) + \deg(V) \geq n$$

$\therefore$  given graph has Hamiltonian Circuit

Hence theorem is proved.

**(B) Define partial order set, chain, lexicographic, Isomorphism and show that the function  $f: A \rightarrow A'$  define by  $f(a) = 2(a)$  is an isomorphism from  $(A, \leq)$  to**

**$(A', \leq)$  where  $A$  is a set of positive integers,  $A'$  is a set of positive even integers.**

**Solution:**

partial order set :

Let  $A$  is a relation on set  $A$ . then relation  $R$  is called partial order. If it is reflexive, antisymmetric and transitive.

If  $R$  is a partial order relation on set  $A$ . then set  $A$  together with partial order relation  $R$  is known as partial order set or partial order set.

Ex. Let  $Z$  be a set of integers " $\leq$ " be a relation on  $Z$ .

$\therefore$  Reflexive property is satisfied.

$$(\because a \leq a \quad \forall a \in Z)$$

Let  $a, b \in Z$

$$a \leq b \text{ and } b \leq a \Rightarrow a = b$$

$\therefore$  Antisymmetric property is satisfied

$$a \leq b \text{ and } b \leq c \Rightarrow a \leq c$$

$\therefore$  Transitive property is satisfied.

$\therefore$  " $\leq$ " is a partial order relation on  $Z$

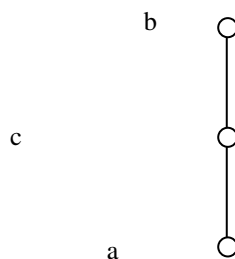
Similarly " $\geq$ " is also a partial order relation on  $Z$ .

Chain order set:

If every pair of element in a poset is comparable then poset  $A$  is called linear order set. Or set  $A$  is chain.

Ex.  $A = \{a, b, c\}$

$$a \leq c, c \leq b$$



This order is in linear or chain.

Hence it is called chain or linear order.

Lexicographic:

Let  $A \times B$  is a cartesian product of two sets  $A$  &  $B$ . we define " $<$ " as follow.

$(a, b) < (a', b')$  if  $a < a'$  or if  $a = a'$  then  $b < b'$

This is used in dictionary.

Hence it is also as dictionary

Ex. Help, help

Help  $<$  help

Isomorphism:

Let  $(A, \leq)$  and  $(A', \leq')$  be posets and let  $f: A \rightarrow A'$  be a one-to-one correspondence between  $f: A$  &  $A'$ ! The function  $f$  is called an Isomorphism from  $A$  to  $A'$

It for any  $a, b \in A, a \leq b$

$\Leftrightarrow f(a) \leq' f(b)$ .

**(Proof left)**

**(c) Define bounded lattice, distributive lattice, complemented lattice, modular lattice and prove that if  $L$  is a bounded distributive lattice then if complement exists, it is unique.**

**Solution:**

**(D) For Boolean polynomial**

$$P(x, y, z) = (x \wedge y) \vee (y \wedge z)$$

**Construct truth table and show the polynomial by logic diagram.**